

ECE-374-B: Lecture 0 - Logistics and Strings/Languages

Lecturer: Nickvash Kani

August 27, 2024

University of Illinois at Urbana-Champaign

Course Administration

Course Policies

See website

Discussion Sessions/Labs

- 50min problem solving session led by TAs
- Two times a week
- Go to your assigned discussion section
- Bring pen and paper!

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Discussion sections will have questions that appear on the homework. If, you skip, you're just making more work for yourself later.

Any questions

Again all policy information should be on course website:

<https://ecealgo.com/fa24/>

Any questions?

Over-arching course questions

High-Level Questions

This course introduces three distinct fields of computer science research:

- Computational complexity.
 - Given infinite time and a certain machine, is it possible to solve a given problem.
- Algorithms
 - Given a deterministic Turing machine, how fast can we solve certain problems.
- Limits of computation.
 - Are there tasks that our computers cannot do and how do we identify these problems?

Why not just focus on Algorithms?

When someone asks you, "How fast can you compute problem X ", they are actually asking:

- Is X solvable using the deterministic Turing machines we have at our disposal?
- If it is solvable, can we find the solution efficiently (in poly-time)?
- If it is solvable but we don't have a poly time solution, what problem(s) is it most similar too?

Course Structure

Course divided into three parts:

- Basic automata theory: finite state machines, regular languages, hint of context free languages/grammars, Turing Machines
- Algorithms and algorithm design techniques
- Undecidability and NP-Completeness, reductions to prove intractability of problems

Week	Topics/Lectures	Mid Lab	Therapy Session	Pr Lab
Aug 23-24	Introduction and overview of the course Introduction to Automata Theory Beri (Lectures 1-4) Pr. (Week 1-4)	Mathematical Induction Beri (Lectures 5-6) Pr. (Week 5-6)	Formal Languages Lectures and exercises Beri (Lectures 7-8) Pr. (Week 7-8)	Formal Languages Beri (Lectures 9-10) Pr. (Week 9-10)
Aug 30 - Sep 1	Finite Automata Beri (Lectures 11-12) Pr. (Week 11-12)	Regular Expressions Beri (Lectures 13-14) Pr. (Week 13-14)	Finite Automata Lectures and exercises Beri (Lectures 15-16) Pr. (Week 15-16)	Finite Automata Beri (Lectures 17-18) Pr. (Week 17-18)
Sep 8-9	Pushdown Automata Beri (Lectures 19-20) Pr. (Week 19-20)	Context-Free Languages Beri (Lectures 21-22) Pr. (Week 21-22)	Pushdown Automata Lectures and exercises Beri (Lectures 23-24) Pr. (Week 23-24)	Pushdown Automata Beri (Lectures 25-26) Pr. (Week 25-26)
Sep 15-16	Context-Free Languages and Grammars Beri (Lectures 27-28) Pr. (Week 27-28)	Context-Free Languages and Grammars Beri (Lectures 29-30) Pr. (Week 29-30)	Turing Machines Lectures and exercises Beri (Lectures 31-32) Pr. (Week 31-32)	Turing Machines Beri (Lectures 33-34) Pr. (Week 33-34)
Sep 22-23	Context-Free Languages and Grammars Beri (Lectures 35-36) Pr. (Week 35-36)	Context-Free Languages and Grammars Beri (Lectures 37-38) Pr. (Week 37-38)	Midweek 3 - Thursday Sep 23 (Week 35-36) Beri (Lectures 39-40) Pr. (Week 39-40)	No interaction
Sep 29-30	Reductions and Normalization Beri (Lectures 41-42) Pr. (Week 41-42)	Reductions and Normalization Beri (Lectures 43-44) Pr. (Week 43-44)	Diagrams and context: Reductions, Normalization Beri (Lectures 45-46) Pr. (Week 45-46)	Diagrams and Context Beri (Lectures 47-48) Pr. (Week 47-48)
Oct 7-8	Reductions and Normalization Beri (Lectures 49-50) Pr. (Week 49-50)	Reductions and Normalization Beri (Lectures 51-52) Pr. (Week 51-52)	Complexity Lectures and exercises Beri (Lectures 53-54) Pr. (Week 53-54)	Complexity Beri (Lectures 55-56) Pr. (Week 55-56)
Oct 14-15	More Complexity Beri (Lectures 57-58) Pr. (Week 57-58)	More Complexity Beri (Lectures 59-60) Pr. (Week 59-60)	Complexity Lectures and exercises Beri (Lectures 61-62) Pr. (Week 61-62)	Complexity Beri (Lectures 63-64) Pr. (Week 63-64)
Oct 21-22	Complexity Beri (Lectures 65-66) Pr. (Week 65-66)	Complexity Beri (Lectures 67-68) Pr. (Week 67-68)	Complexity Lectures and exercises Beri (Lectures 69-70) Pr. (Week 69-70)	Complexity Beri (Lectures 71-72) Pr. (Week 71-72)
Oct 29-30	Complexity Beri (Lectures 73-74) Pr. (Week 73-74)	Complexity Beri (Lectures 75-76) Pr. (Week 75-76)	Midweek 4 - Friday Sep 29 (Week 73-74) Beri (Lectures 77-78) Pr. (Week 77-78)	Complexity Beri (Lectures 79-80) Pr. (Week 79-80)
Nov 4-5	Complexity Beri (Lectures 81-82) Pr. (Week 81-82)	Complexity Beri (Lectures 83-84) Pr. (Week 83-84)	Complexity Lectures and exercises Beri (Lectures 85-86) Pr. (Week 85-86)	Complexity Beri (Lectures 87-88) Pr. (Week 87-88)
Nov 11-12	Complexity Beri (Lectures 89-90) Pr. (Week 89-90)	Complexity Beri (Lectures 91-92) Pr. (Week 91-92)	Complexity Lectures and exercises Beri (Lectures 93-94) Pr. (Week 93-94)	Complexity Beri (Lectures 95-96) Pr. (Week 95-96)
Nov 18-19	Complexity Beri (Lectures 97-98) Pr. (Week 97-98)	Complexity Beri (Lectures 99-100) Pr. (Week 99-100)	Complexity Lectures and exercises Beri (Lectures 101-102) Pr. (Week 101-102)	Complexity Beri (Lectures 103-104) Pr. (Week 103-104)
Nov 25-26	Complexity Beri (Lectures 105-106) Pr. (Week 105-106)	Complexity Beri (Lectures 107-108) Pr. (Week 107-108)	Complexity Lectures and exercises Beri (Lectures 109-110) Pr. (Week 109-110)	Complexity Beri (Lectures 111-112) Pr. (Week 111-112)
Dec 2-3	Complexity Beri (Lectures 113-114) Pr. (Week 113-114)	Complexity Beri (Lectures 115-116) Pr. (Week 115-116)	Complexity Lectures and exercises Beri (Lectures 117-118) Pr. (Week 117-118)	Complexity Beri (Lectures 119-120) Pr. (Week 119-120)
Dec 9-10	Complexity Beri (Lectures 121-122) Pr. (Week 121-122)	Complexity Beri (Lectures 123-124) Pr. (Week 123-124)	Complexity Lectures and exercises Beri (Lectures 125-126) Pr. (Week 125-126)	Complexity Beri (Lectures 127-128) Pr. (Week 127-128)

Goals

- Algorithmic thinking
- Learn/remember some basic tricks, algorithms, problems, ideas
- Understand/appreciate limits of computation (intractability)
- Appreciate the importance of algorithms in computer science and beyond (engineering, mathematics, natural sciences, social sciences, ...)

Formal languages and complexity (The Blue Weeks!)

Why Languages?

First 5 weeks devoted to language theory.

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But why study languages?

Multiplying Numbers

Consider the following problem:

Problem Given two n -digit numbers x and y , compute their product.

Grade School Multiplication

Compute “partial product” by multiplying each digit of y with x and adding the partial products.

$$\begin{array}{r} 3141 \\ \times 2718 \\ \hline 25128 \\ 3141 \\ 21987 \\ 6282 \\ \hline 8537238 \end{array}$$

Time analysis of grade school multiplication

- Each partial product: $\Theta(n)$ time
- Number of partial products: $\leq n$
- Adding partial products: n additions each $\Theta(n)$ (Why?)
- Total time: $\Theta(n^2)$
- Is there a faster way?

Fast Multiplication

- $O(n^{1.58})$ time [Karatsuba 1960] disproving Kolmogorov's belief that $\Omega(n^2)$ is best possible
- $O(n \log n \log \log n)$ [Schönhage-Strassen 1971].
Conjecture: $O(n \log n)$ time possible
- $O(n \log n \cdot 2^{O(\log^* n)})$ time [Furer 2008]
- $O(n \log n)$ [Harvey-van der Hoeven 2019]

Can we achieve $O(n)$? No lower bound beyond trivial one!

Equivalent Complexity

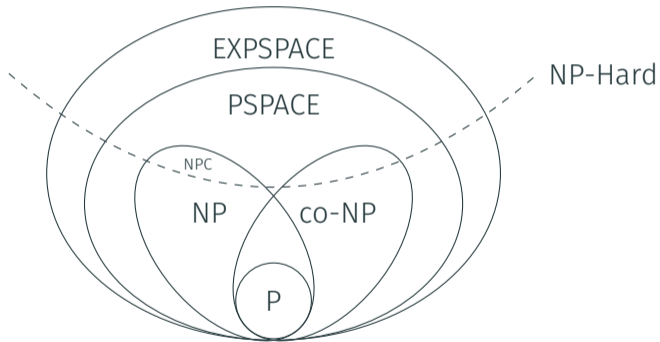
Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort?

Equivalent Complexity

Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort? How do we compare? The two problems have:

- Different inputs (two numbers vs n-element array)
- Different outputs (a number vs n-element array)
- Different entropy characteristics (from a information theory perspective)

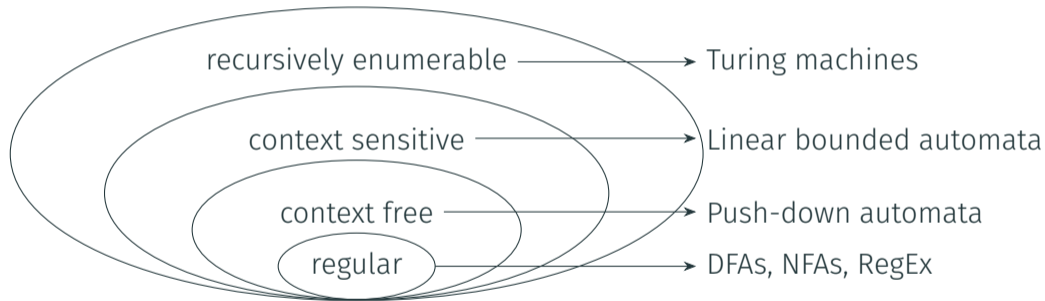
An algorithm has a runtime complexity.



Languages, Problems and Algorithms ... oh my! III

A problem has a complexity class!

Recognized by:



Problems do not have run-time since a problem \neq the algorithm used to solve it.

Complexity classes are defined differently.

How do we compare problems? What if we just want to know if a problem is "computable"

Definition

1. An **algorithm** is a step-by-step way to solve a problem.
2. A **problem** is some question that we'd like answered given some input. It should be a decision problem of the form "Does a given input fulfill property X ."
3. A **Language** is a set of strings. Given an alphabet, Σ a language is a subset of Σ^*

Definition

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$L = \{w \mid w \text{ is the encoding of an input } y \text{ to problem } X \text{ and the answer to input } y \text{ for a problem } X \text{ is "YES"} \}$

A decision problem X is "YES" is the string is in the language.

Language of multiplication

How do we define the multiplication problem as a language?

Define L as language where inputs are separated by comma and output is separated by $|$.

Machine accepts a $x*y=z$ if " $x*y|z$ " is in L . Rejects otherwise.

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$$L_{MULT2} = \left\{ \begin{array}{lll} 1 \times 1|1, & 1 \times 2|2, & 1 \times 3|3, \dots \\ 2 \times 1|2, & 2 \times 2|4, & 2 \times 3|6, \dots \\ \vdots & \vdots & \vdots \\ n \times 1|n, & n \times 2|2n, & n \times 3|3n, \dots \end{array} \right\} \quad (1)$$

Language of sorting

We do the same thing for sorting.

Define L as language where inputs are separated by comma and output is separated by |.

Machine accepts a $[i_1, i_2, \dots] = \text{sort}(\{i_1, i_2, \dots\})$ if " $x[]|z[]$ " is in L. Rejects otherwise.

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$$L_{\text{Sort2}} = \left\{ \begin{array}{ccc} 1, 1|1, 1 & 1, 2|1, 2 & 1, 3|1, 3, \dots \\ 2, 1|1, 2, & 2, 2|2, 2, & 2, 3|2, 3, \dots \\ \vdots & \vdots & \vdots \\ n, 1|1, n, & n, 2|2, n, & n, 3|3, n, \dots \end{array} \right\} \quad (2)$$

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If the same type of machine can recognize both languages, then that gives us an upperbound to their hardness.

How do we formulate languages?

Strings

Alphabet

An **alphabet** is a **finite** set of symbols.

Examples of alphabets:

- $\Sigma = \{0, 1\}$,
- $\Sigma = \{a, b, c, \dots, z\}$,
- ASCII.
- UTF8.
- $\Sigma = \{\langle(w)forward\rangle, \langle(a)strafe\ left\rangle, \langle(s)back\rangle, \langle(d)strafe\ right\rangle\}$

String Definition

Definition

1. A **string/word** over Σ is a **finite sequence** of symbols over Σ . For example, '0101001', '*string*', ' \langle moveback \rangle \langle rotate90 \rangle '
2. $x \cdot y \equiv xy$ is the concatenation of two strings
3. The **length** of a string w (denoted by $|w|$) is the number of symbols in w . For example, $|101| = 3$, $|\epsilon| = 0$
4. For integer $n \geq 0$, Σ^n is set of all strings over Σ of length n . Σ^* is the set of all strings over Σ .
5. Σ^* set of all strings of all lengths including empty string.

Question: $\{ '0', '1' \}^* =$

- ϵ is a **string** containing no symbols. It is not a set
- $\{\epsilon\}$ is a **set** containing one string: the empty string. It is a set, not a string.
- \emptyset is the **empty set**. It contains no strings.

Question: What is $\{\emptyset\}$

Concatenation and properties

- If x and y are strings then xy denotes their concatenation.
- **Concatenation** defined recursively :
 - $xy = y$ if $x = \epsilon$
 - $xy = a(wy)$ if $x = aw$
- xy sometimes written as $x \bullet y$.
- concatenation is **associative**: $(uv)w = u(vw)$ hence write $uvw \equiv (uv)w = u(vw)$
- **not** commutative: uv not necessarily equal to vu
- The identity element is the empty string ϵ :

$$\epsilon u = u \epsilon = u.$$

Definition

v is **substring** of $w \iff$ there exist strings x, y such that $w = xvy$.

- If $x = \epsilon$ then v is a **prefix** of w
- If $y = \epsilon$ then v is a **suffix** of w

Subsequence

A subsequence of a string $w[1..n]$ is either a subsequence of $w[2..n]$ or $w[1]$ followed by a subsequence of $w[2..n]$.

Example

EE37 is a subsequence of *ECE374B*

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Question: How many sub-sequences are there in a string $|w| = 6$?

Definition

If w is a string then w^n is defined inductively as follows:

$$w^n = \epsilon \text{ if } n = 0$$

$$w^n = ww^{n-1} \text{ if } n > 0$$

Question: $(ha)^3 =$.

Rapid-fire questions -strings

Answer the following questions taking $\Sigma = \{0, 1\}$.

1. What is Σ^0 ?
2. How many elements are there in Σ^n ?
3. If $|u| = 2$ and $|v| = 3$ then what is $|u \cdot v|$?
4. Let u be an arbitrary string in Σ^* . What is ϵu ? What is $u \epsilon$?

Languages

Definition

A **language** L is a set of strings over Σ . In other words $L \subseteq \Sigma^*$.

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Standard set operations apply to languages.

- For languages A, B the **concatenation** of A, B is $AB = \{xy \mid x \in A, y \in B\}$.
- For languages A, B , their **union** is $A \cup B$, **intersection** is $A \cap B$, and **difference** is $A \setminus B$ (also written as $A - B$).
- For language $A \subseteq \Sigma^*$ the **complement** of A is $\bar{A} = \Sigma^* \setminus A$.

Set Concatenation

Definition

Given two sets X and Y of strings (over some common alphabet Σ) the **concatenation** of X and Y is

$$XY = \{xy \mid x \in X, y \in Y\} \quad (3)$$

Question: $X = \{ECE, CS, \}$, $Y = \{340, 374\} \implies$
 $XY = .$

Definition

1. Σ^n is the set of all strings of length n . Defined inductively:

$$\Sigma^n = \{\epsilon\} \text{ if } n = 0$$

$$\Sigma^n = \Sigma\Sigma^{n-1} \text{ if } n > 0$$

2. $\Sigma^* = \cup_{n \geq 0} \Sigma^n$ is the set of all finite length strings
3. $\Sigma^+ = \cup_{n \geq 1} \Sigma^n$ is the set of non-empty strings.

Definition

A **language** L is a set of strings over Σ . In other words $L \subseteq \Sigma^*$.

Question: Does Σ^* have strings of infinite length?

Problem

Consider languages over $\Sigma = \{0, 1\}$.

1. What is \emptyset^0 ?
2. If $|L| = 2$, then what is $|L^4|$?
3. What is \emptyset^* , $\{\epsilon\}^*$?
4. For what L is L^* finite?
5. What is \emptyset^+ ?
6. What is $\{\epsilon\}^+$?

Terminology Review

Let's review what we learned.

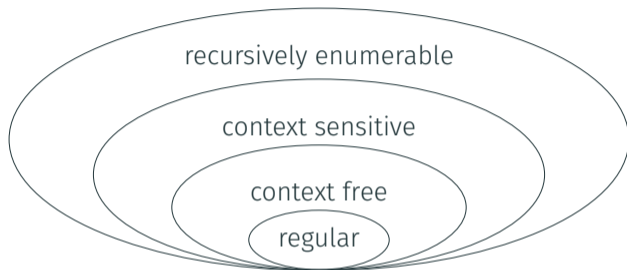
- A **character**(a, b, c, x) is a unit of information represented by a symbol: (letters, digits, whitespace)
- A **alphabet**(Σ) is a set of characters
- A **string**(w) is a sequence of characters
- A **language**(A, B, C, L) is a set of strings

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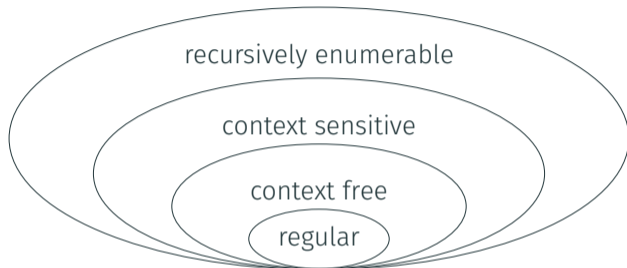
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- A **string**(w) is a sequence of characters
- A **language**(A, B, C, L) is a set of strings
- A **grammar**(G) is a set of rules that defines the strings that belong to a language

Languages: easiest, easy, hard, really hard, reallyⁿ hard



- Regular languages.
 - Regular expressions.
 - DFA: Deterministic finite automata.
 - NFA: Non-deterministic finite automata.
- Context free languages (stack).
- Turing machines: Decidable languages.
- TM Undecidable/unrecognizable languages (halting theorem).

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- Regular languages.
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