

1 DFAs

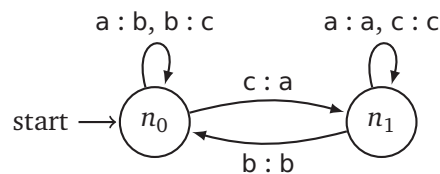
Describe deterministic finite-state automata that accept each of the following languages over the alphabet $\Sigma = \{0, 1\}$. Describe briefly what each state in your DFAs *means*.

Either drawings or formal descriptions are acceptable, as long as the states Q , the start state s , the accept states A , and the transition function δ are all clear. Try to keep the number of states small.

1. All strings containing the substring **000**.
2. All strings *not* containing the substring **000**.
3. Every string except **000**. [*Hint: Don't try to be clever.*]
4. All strings in which the number of **0**s is even **and** the number of **1**s is *not* divisible by 3.
5. All strings in which the number of **0**s is even **or** the number of **1**s is *not* divisible by 3.
6. Given DFAs M_1 and M_2 , all strings in $\overline{L(M_1)} \oplus L(M_2)$.
Recall that for two sets A and B , their symmetric distance $A \oplus B$ is the set of elements in either A or B , but not both.

2 Other types of automata

1. A *finite-state transducer* (FST) is a type of deterministic finite automaton whose output is a string instead of just *accept* or *reject*. The following is the state diagram of finite state transducer FST_0 .



Each transition of an FST is labeled at least an input symbol and an output symbol, separated by a colon (:). There can also be multiple input-output pairs for each transitions, separated by a comma (,). For instance, the transition from n_0 to itself can either take a or b as an input, and outputs b or c respectively.

When an FST computes on an input string $s := \overline{s_0s_1\dots s_{n-1}}$ of length n , it takes the input symbols s_0, s_1, \dots, s_{n-1} one by one, starting from the starting state, and produces corresponding output symbols. For instance, the input string $abccba$ produces the output string $bcacbb$, while $cbaabc$ produces $abbbca$.

- (a) Each of the following strings is the input of FST_0 . Give the sequence of states entered and the output produced.
 - aaca
 - cbbc
 - bcba

- acbbca
- (b) Assume that FST's have an input alphabet Σ and an output alphabet Γ , give a formal definition of this type of model and its computation. (Hint: An FST is a 5-tuple with no accepting states. Its transition function is of the form $\delta : Q \times \Sigma \rightarrow Q \times \Gamma$.)
 - (c) Give a formal description of FST_0 .
 - (d) Give a state diagram of an FST with the following behavior. Its input and output alphabets are $\{T, F\}$. Its output string is inverted on the positions with indices divisible by 3 and is identical on all the other positions. For instance, on an input TFTTFTFT it should output FFTFFTTT.

Work on these later:

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7. All strings w such that *in every prefix of w* , the number of **0**s and **1**s differ by at most 1.
8. All strings containing at least two **0**s and at least one **1**.
9. All strings w such that *in every prefix of w* , the number of **0**s and **1**s differ by at most 2.
- *10. All strings in which the substring **000** appears an even number of times.
(For example, **0001000** and **0000** are in this language, but **00000** is not.)
11. All strings that are **both** the binary representation of an integer divisible by 3 **and** the ternary (base-3) representation of an integer divisible by 4.

For example, the string **1100** is an element of this language, because it represents $2^3 + 2^2 = 12$ in binary and $3^3 + 3^2 = 36$ in ternary.
- *12. All strings w such that $F_{\#(\mathbf{10}, w)} \bmod 10 = 4$, where $\#(\mathbf{10}, w)$ denotes the number of times **10** appears as a substring of w , and F_n is the n th Fibonacci number:

$$F_n = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ F_{n-1} + F_{n-2} & \text{otherwise} \end{cases}$$