Recall fooling sets and distinguishability. Two strings $x, y \in \Sigma^*$ are suffix distinguishable with respect to a given language L if there is a string z such that exactly one of xz and yz is in L. This means that any DFA that accepts L must necessarily take x and y to different states from its start state. A set of strings F is a fooling set for L if *any* pair of strings $x, y \in F, x \neq y$ are distinguisable. This means that any DFA for L requires at least $|F|$ states. To prove non-regularity of a language L you need to find an infinite fooling set F for L. Given a language L try to find a constant size fooling set first and then prove that one of size n exists for any given n which is basically the same as finding an infinite fooling set.

Note that another method to prove non-regularity is via *reductions*. Suppose you want to prove that L is non-regular. You can do regularity preserving operations on L to obtain a language L' which you already know is non-regular. Then L must not have been regular. For instance if $\bar{L}$ is not regular then L is also not regular. You will see an example in Problem 4 below.

Prove that each of the following languages is ***not*** regular.

1. $\{\mathtt{0}^{2n}\mathtt{1}^n \mid n \geq 0\}$

2. $\{\mathtt{0}^m\mathtt{1}^n \mid m \neq 2n\}$

3. $\left\{\mathtt{0}^{2^n} \;\middle|\; n \geq 0\right\}$

4. Strings over $\{\mathtt{0}, \mathtt{1}\}$ where the number of $\mathtt{0}$s is exactly twice the number of $\mathtt{1}$s.
   - Describe an infinite fooling set for the language.
   - Use closure properties. What is language if you intersect the given language with $\mathtt{0}^*\mathtt{1}^*$?

5. Strings of properly nested parentheses **()**, brackets **[]**, and braces **{}**. For example, the string **([]){}** is in this language, but the string **([)]** is not, because the left and right delimiters don't match.
   - Describe an infinite fooling set for the language.
   - Use closure properties.

6. w, such that $|w| = \lceil k\sqrt{k} \rceil$, for some natural number k.

   Hint: since this one is more difficult, we'll even give you a fooling set that works: try $F = \{\mathtt{0}^{m^6} | m \geq 1\}$. We'll also provide a bound that can help: the difference between consecutive strings in the language, $\lceil (k+1)^{1.5} \rceil - \lceil k^{1.5} \rceil$, is bounded above and below as follows
   $$1.5\sqrt{k} - 1 \leq \lceil (k+1)^{1.5} \rceil - \lceil k^{1.5} \rceil \leq 1.5\sqrt{k} + 3$$
   All that's left is you need to carefully prove that F is a fooling set for L.

7. Strings of the form $w_1\mathtt{\#}w_2\mathtt{\#}\cdots\mathtt{\#}w_n$ for some $n \geq 2$, where each substring w_i is a string in $\{\mathtt{0}, \mathtt{1}\}^*$, and some pair of substrings w_i and w_j are equal.

Work on these later:

7. $\left\{\mathtt{0}^{n^2} \mid n \geq 0\right\}$

8. $\{w \in (\mathtt{0} + \mathtt{1})^* \mid w$ is the binary representation of a perfect square$\}$