CEECE-374-B: Lecture 0 - Logistics andStrings/Languages

Lecturer: Nickvash Kani

August 22, 2023

University of Illinois at Urbana-Champaign

Course Administration

Instructional Staff

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- Contacting us: Use <u>private notes</u> on Piazza to reach course staff. Direct email only for sensitive or confidential information.

This semester, the two sections will be run completely independently.

- Different lectures.
- Different homeworks, quizzes, exams.
- Different grading policies.

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Section B will be in-person only. Recordings will be attempted but not guaranteed.

- Webpage: General information, announcements, homeworks, quizzes, course policies https://ecealgo.com
- Submission(Gradescope): Written homework submission and grading, regrade requests. Exams will be uploaded there as well.
- Communication(Piazza): Announcements, online questions and discussion, contacting course staff (via private notes)
- Gradebook (Canvas): Announcements, online questions and discussion, contacting course staff (via private notes)

See course webpage for links

Important: check Piazza/course web page at least once each day

Discussion Sessions/Labs

- 50min problem solving session led by TAs
- Two times a week
- Go to your assigned discussion section
- Bring pen and paper!

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Discussion sections will have questions that appear on the homework. If, you skip, you're just making more work for yourself later.

Again all policy information should be on course website: https://ecealgo.com

Any questions?

Over-arching course questions

This course introduces three distinct fields of computer science research:

- Computational complexity.
 - Given infinite time and a certain machine, is it possible to solve a given problem.
- Algorithms
 - Given a deterministic Turing machine, how fast can we solve certain problems.
- Limits of computation.
 - Are there tasks that our computers cannot do and how do we identify these problems?

When someone asks you, "How fast can you compute problem *X*", they are actually asking:

- Is X solvable using the deterministic Turing machines we have at our disposal?
- If it is solvable, can we find the solution efficiently (in poly-time)?
- If it is solvable but we don't have a poly time solution, what problem(s) is it most similar too?

Course divided into three parts:

- Basic automata theory: finite state machines, regular languages, hint of context free languages/grammars, Turing Machines
- Algorithms and algorithm design techniques
- Undecidability and NP-Completeness, reductions to prove intractability of problems

Week	Tuesday Lecture	Wed Lab	Thursday Lecture	Fri Lab
Aug 23-26	Adminis trivia and course goals Introduction and history 🕹 : strings 🕹 (Sarie's Videos, Lec. 1 e) [B: slides, scribbles, video]	String induction & [Jeff's induction notes . Chandra's induction notes of [solutions & [recording]	Languages and regular expressions 🛓 [Sariet's Videos, Lec. 2 e] [B: slides, scribble]	Regular expressions
Aug 30 - Sep 2	DFAs: intuition, definitions, closure properties & Automata Tutor et , JELAP et , Mahesh's DFA notes et , Sarki's Videos, Lec 3 et] [8: sildes, snibble]	DFA construction 4 (solutions 4) (recording)	Non-Determinism, NFAs J [Sarief's Videos, Lec.4 et] [B: sildes, scribble]	Language transformations & [solutions &] [recording]
Sep 6-9	Equivalence of DFAs, NFAs, and regular expressions & [Sariet's Videos, Lec 5 of] [B: slides, scribble]	Regex to NFA to DFA (to Regex) & (solutions &) (recording)	Fooling Sets and Proving Non-Regularity Undershis DFA notes σ , Fall 2015 TAs' Fooling Sets Notes σ , Sarief's Videos, Lec é σ) [8: slides, scribble]	Proving Non- Regularity & [solutions & [recording]
Sep 13-16	Context-free languages and grammars [Sariel's Videos, Lec 7 #] [B: slides, scribble]	Context-free grammars 🕹 [solutions 🕁 [recording]	Turing machines: history, formal definitions, examples, variations [Sariet's Videos, Lec 8 of] [8: slides, scribble]	Turing Machines. & [solutions &][recording]
Sep 20-23	Universal Turing machines 4 [Sariel's Videos, Lec 8 or] [B: slides, scribble]	Midterm 1 Review (effectively office hours)	<u>Midterm 1</u> - Thursday Sep 22:12:30-13:45 [skill-set ↓, fodder ↓, cheat sheet ↓]	No Instruction
Sep 27-30	Reductions & Recursion J [Sarief's Videoc, Lec 10 e , Notes on Solving Recurrences o] [B: slides, scribbles]	Binary search 🕹 (solutions 🛓 [(recording]	Divide and conquer: Selection, Karatsuba J. (Sarie's Videos, Lec 11 of] (B: slides, scribbles)	Divide and Conquer
Oct 4-7	Backtracking, 🛓 (Sariefs Videos, Lec 12 et) [B: slides, scribbles]	Backtracking 🛓 (solutions 🛓 (recording)	Dynamic programming, 🛓 (Sariet's Videos, Lec 13 et) [B: slides, scribbles]	Dynamic programming \downarrow (solutions \downarrow [recording]
Oct 11-14	More Dynamic programming [Sariel's Videos, Lec 14 et] [B: slides, scribbles]	More Dynamic programming [solutions [recording]	Graphs, Basic Search & (Chandra's Graph notes & , Sariel's Videos, Lec 15 &] (B: slides, scribbles]	Even more DP 🞍 [solutions 🛓](recording)
Oct 18-21	Directed Graphs. DFS, DAGs and Topological Sort J. [Chandra's Graph notes a . Sariel's Videos, Lec 16 all [ft:sildes, scribbles]	Graph Modeling 🕹 (solutions 🛓 [(recording]	Shortest Paths: BFS and Dijkstra. J. [Chandra's Graph notes e , Sariet's Videos, Lec 17 e1 [B: slides, scribbles]	Shortest Paths 🕹 Isolutions 🛓 [(recording]
Oct 25-28	Bellman-Ford, Dynamic Programming on DAGs ↓ (Chandra's Graph notes σ , Sariel's Videos, Lec 18 c ⁺] (B: slides, scribbles]	More Shortest Paths	MST Algorithms. J.	MST ↓ [solutions ↓](recording]
Nov 1-4	Midterm 2 (Recursion/DP/Graph Algorithms) - Tuesday, Nov 1 12:30-13:45 [skill-set ↓, fodder ↓, cheat sheet ↓]	No Instruction	Reductions & (Sariel's Videos, Lec 21 et) (B: slides, scribbles)	Reductions 🛓 (solutions 🛓 (recording)
Nov 8-11	SAT, NP and NP-Hardness # [Sarief's Videos, Lec 22-24 of] [B: slides, scribbles]	NP-hardness reductions & [solutions & [recording]	More NP-Hardness [Sariel's Visicos, Lec. 23-24 of] [B: slides, scribbles]	More NP-Hardness ↓ [solutions ↓ [recording]
Nov 15-18	Undecidability 1 J [Sariel's Videos, Lec. 9 tř.] [B: slides, scribbles]	Undecidability reductions 🛓 (solutions 🞍 [recording]	Undecidability 2 [Sariel's Videos, Lec. 9 et] [8: slides, scribbles]	No instruction
Fall Break (Nov 19-27). Have fun.				
Nov 29 - Dec 2	Optional review for Midterm 3 [B: slides, scribbles]	Optional Review for midterm 3 (effectively office hours)	<u>Midterm 3</u> (Reductions/P- NP/Decidability) - Thursday, Dec 1 15:30-16:45 [ان غ اناً (ان غان), fodder ان , <u>cheat sheet</u> ()	
Dec 6-9	Wrap-up, closing remarks Optional review for Final Exam [B: slides, scribbles]	Optional Review for final exam	Reading Day ICES Forms Due	
Final Exam – TBD [Skill set 🕹 , cheat sheet 🛃]				

Goals

Algorithmic thinking

- Learn/remember some basic tricks, algorithms, problems, ideas
- Understand/appreciate limits of computation (intractability)
- Appreciate the importance of algorithms in computer science and beyond (engineering, mathematics, natural sciences, social sciences, ...)

Formal languages and complexity (The Blue Weeks!)

First 5 weeks devoted to language theory.

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But why study languages?

Consider the following problem:

Problem Given two *n*-digit numbers *x* and *y*, compute their product.

Grade School Multiplication

Compute "partial product" by multiplying each digit of y with x and adding the partial products.





Time analysis of grade school multiplication

- Each partial product: $\Theta(n)$ time
- Number of partial products: $\leq n$
- Adding partial products: n additions each $\Theta(n)$ (Why?)
- Total time: $\Theta(n^2)$
- Is there a faster way?

Fast Multiplication

- $O(n^{1.58})$ time [Karatsuba 1960] disproving Kolmogorov's belief that $\Omega(n^2)$ is best possible
- O(n log n log log n) [Schonhage-Strassen 1971].
 Conjecture: O(n log n) time possible
- $O(n \log n \cdot 2^{O(\log^* n)})$ time [Furer 2008]
- $O(n \log n)$ [Harvey-van der Hoeven 2019]

Can we achieve O(n)? No lower bound beyond trivial one!

Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort?

Does this mean multiplication is as complex as another problem that has a *O*(*n* log *n*) algorithm like sorting/QuickSort? How do we compare? The two problems have:

- Different inputs (two numbers vs n-element array)
- Different outputs (a number vs n-element array)
- Different entropy characteristics (from a information theory perspective)

An algorithm has a runtime complexity.



Languages, Problems and Algorithms ... oh my! III

A problem has a complexity class!



Problems do not have run-time since a problem \neq the algorithm used to solve it. *Complexity classes are defined differently.*

How do we compare problems? What if we just want to know if a problem is "computable".

Algorithms, Problems and Languages ... oh my! I

Definition

- 1. An algorithm is a step-by-step way to solve a problem.
- A problem is some question that we'd like answered given some input. It should be a decision problem of the form "Does a given input fulfill property X."
- 3. A Language is a set of strings. Given a alphabet, Σ a language is a subset of Σ^{\ast}

Algorithms, Problems and Languages ... oh my! I

Definition

- 1. An algorithm is a step-by-step way to solve a problem.
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- A Language is a set of strings. Given a alphabet, Σ a language is a subset of Σ* A language is a formal realization of this problem. For problem X, the corresponding language is:

L = {w | w is the encoding of an input y to problem X and the answer to input y for a problem X is "YES" } A decision problem X is "YES" is the string is in the language. How do we define the multiplication problem as a language?

Define L as language where inputs are separated by comma and output is separated by |.

Machine accepts a x*y=z if "x*y|z" is in L. Rejects otherwise.

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$$L_{MULT2} = \begin{cases} 1 \times 1|1, & 1 \times 2|2, & 1 \times 3|3, \dots \\ 2 \times 1|2, & 2 \times 2|4, & 2 \times 3|6, \dots \\ \vdots & \vdots & \vdots \\ n \times 1|n, & n \times 2|2n, & n \times 3|3n, \dots \end{cases}$$
(1)

We do the same thing for sorting.

Define L as language where inputs are separated by comma and output is separated by |.

Machine accepts a $[i_1, i_2, \ldots] = sort(\{i_1, i_2, \ldots\})$ if "x[]|z[]" is in L. Rejects otherwise.

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$$L_{Sort2} = \begin{cases} 1, 1|1, 1 & 1, 2|1, 2 & 1, 3|1, 3, \dots \\ 2, 1|1, 2, & 2, 2|2, 2, & 2, 3|2, 3, \dots \\ \vdots & \vdots & \vdots \\ n, 1|1, n, & n, 2|2, n, & n, 3|3, n, \dots \end{cases}$$

(2)

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(2)

If the same type of machine can recognize both languages, then that gives us an upperbound top their hardness. How do we formulate languages?

Strings

An alphabet is a **finite** set of symbols.

Examples of alphabets:

- $\Sigma = \{0, 1\},$
- $\Sigma = \{a, b, c, \dots, z\}$,
- ASCII.
- UTF8.
- $\Sigma = \{\langle (w) \text{forward} \rangle, \langle (a) \text{strafe left} \rangle, \langle (s) \text{back} \rangle, \langle (d) \text{strafe right} \rangle \}$

Definition

- A string/word over Σ is a finite sequence of symbols over
 Σ. For example, '0101001', 'string', '(moveback)(rotate90)'
- 2. $x \cdot y \equiv xy$ is the concatenation of two strings
- 3. The length of a string w (denoted by |w|) is the number of symbols in w. For example, |101| = 3, $|\epsilon| = 0$
- 4. For integer $n \ge 0$, Σ^n is set of all strings over Σ of length n. Σ^* is the set of all strings over Σ . $Z^* = \underbrace{\mathcal{E}} \cdot \underbrace{\mathcal{E}} \cdot$
- 5. Σ^* set of all strings of all lengths including empty string.

- + ϵ is a string containing no symbols. It is not a set
- \cdot ({ ϵ } is a set containing one string: the empty string. It is a set, not a string.

 $\stackrel{!}{=} \emptyset$ is the empty set. It contains no strings. Question: What is $\{\emptyset\} = \emptyset$

$$\phi^2 = O$$

Concatenation and properties

- If x and y are strings then xy denotes their concatenation.
- Concatenation defined recursively :
 - xy = y if $x = \epsilon$
 - xy = a(wy) if x = aw
- *xy* sometimes written as *x y*.
- concatenation is associative: (uv)w = u(vw) hence write $uvw \equiv (uv)w = u(vw)$
- **not** commutative: *uv* not necessarily equal to *vu*
- The <u>identity</u> element is the empty string ϵ :

$$\epsilon U = U\epsilon = U.$$

Definition v is substring of $w \iff$ there exist strings x, y such that w = xvy.

- If $x = \epsilon$ then v is a prefix of w
- If $y = \epsilon$ then v is a suffix of w

w = vyw = xw A subsequence of a string w[1...n] is either a subsequence of w[2...n] or w[1] followed by a subsequence of w[2...n].

Example *EE37* is a subsequence of *ECE374B* A subsequence of a string w[1...n] is either a subsequence of w[2...n] or w[1] followed by a subsequence of w[2...n].

Example *EE37* is a subsequence of *ECE374B*

Question: How many sub-sequences are there in a string |w| = 6? 2^6



Definition If w is a string then w^n is defined inductively as follows: $w^n = \epsilon \text{ if } n = 0$ $w \cdot w \cdot \ldots \cdot w$ $w^n = ww^{n-1} \text{ if } n > 0$

Question: $(ha)^3 = .$ ho ho ho

Answer the following questions taking $\Sigma = \{0, 1\}$.

- 1. What is Σ^0 ?
- 2. How many elements are there in Σ^n ? 2
- 3. If |u| = 2 and |v| = 3 then what is $|u \cdot v|$?
- 4. Let *u* be an arbitrary string in Σ^* . What is ϵu ? What is $u\epsilon$?
 - U

• C 0,1,00,01 2 10, 11,000 Languages E = all the string where /1 1w1 = 11

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Definition $\mathcal{A}_{\text{language } L}$ is a set of strings over Σ^{*} . In other words $L \subseteq \Sigma^{*}$.

Definition

A language *L* is a set of strings over Σ . In other words $L \subseteq \Sigma^*$.

Standard set operations apply to languages.

- For languages A, B the concatenation of A, B is $AB = \{xy \mid x \in A, y \in B\}.$
- For languages A, B, their union is $A \cup B$, intersection is $A \cap B$, and difference is $A \setminus B$ (also written as A B).
- For language $A \subseteq \Sigma^*$ the complement of A is $\overline{A} = \Sigma^* \setminus A$.

Definition Given two sets X and Y of strings (over some common alphabet Σ) the concatenation of X and Y is

$$XY = \{xy \mid x \in X, y \in Y\}$$
(3)

Question: $X = \{ECE, CS_{y}\}, Y = \{340, 374\} \implies$ XY = . ECE 374, CS 374, Z ECE 340, CS 340

Σ^* and languages

Definition

1. Σ^n is the set of all strings of length *n*. Defined inductively:

- length II. Define $\{\xi\}$ $\{\xi\}$ $\Sigma^n = {\epsilon}$ if n = 0 $\Sigma^n = \Sigma \Sigma^{n-1}$ if n > 02. $\Sigma^* = \bigcup_{n>0} \Sigma^n$ is the set of all finite length strings
- 3. $\Sigma^+ = \bigcup_{n>1} \Sigma^n$ is the set of non-empty strings.

Definition A language L is a set of strings over Σ . In other words $L \subseteq \Sigma^*$.

Question: Does Σ^* have strings of infinite length?

Rapid-Fire questions - Languages

Problem Consider languages over $\Sigma = \{0, 1\}$.

- 1. What is \emptyset^0 ?
- 2. If |L| = 2, then what is $|L^4|$?
- 3. What is \emptyset^* , $\{\epsilon\}^*$?
- 4. For what L is L* finite?
- 5. What is \emptyset^+ ?
- 6. What is $\{\epsilon\}^+$?

Let's review what we learned.

- A character(*a*, *b*, *c*, *x*) is a unit of information represented by a symbol: (letters, digits, whitespace)
- A $alphabet(\Sigma)$ is a set of characters
- A string(w) is a sequence of characters
- A language(A, B, C, L) is a set of strings

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- A language(A, B, C, L) is a set of strings
- A grammar(G) is a set of rules that defines the strings that belong to a language

Languages: easiest, easy, hard, really hard, reallyⁿ hard



- Regular languages.
 - Regular expressions.
 - DFA: Deterministic finite automata.
 - NFA: Non-deterministic finite automata.
 - Languages that are not regular.
- Context free languages (stack).
- Turing machines: Decidable languages.
- TM Undecidable/unrecognizable languages (halting theorem).

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- Regular languages.
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