ECE-374-B: Lecture 0 - Logistics and Strings/Languages

Lecturer: Nickvash Kani
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University of Illinois at Urbana-Champaign

Course Administration

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## Section A vs B

This semester, the two sections will be run completely independently.

- Different lectures.
- Different homeworks, quizzes, exams.
- Different grading policies.


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Section B will be in-person only. Recordings will be attempted but not guaranteed.

## Online resources

- Webpage: General information, announcements, homeworks, quizzes, course policies https://ecealgo.com
- Submission(Gradescope): Written homework submission and grading, regrade requests. Exams wil be uploaded there as well.
- Communication(Piazza): Announcements, online questions and discussion, contacting course staff (via private notes)
- Gradebook (Canvas): Announcements, online questions and discussion, contacting course staff (via private notes)

See course webpage for links
Important: check Piazza/course web page at least once each day

## Discussion Sessions/Labs

- 50min problem solving session led by TAs
- Two times a week
- Go to your assigned discussion section
- Bring pen and paper!


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Discussion sections will have questions that appear on the homework. If, you skip, you're just making more work for yourself later.

## Any questions

Again all policy information should be on course website: https://ecealgo.com

Any questions?

Over-arching course questions

## High-Level Questions

This course introduces three distinct fields of computer science research:

- Computational complexity.
- Given infinite time and a certain machine, is it possible to solve a given problem.
- Algorithms
- Given a deterministic Turing machine, how fast can we solve certain problems.
- Limits of computation.
- Are there tasks that our computers cannot do and how do we identify these problems?


## Why not just focus on Algorithms?

When someone asks you, "How fast can you compute problem $X^{\prime \prime}$, they are actually asking:

- Is X solvable using the deterministic Turing machines we have at our disposal?
- If it is solvable, can we find the solution efficiently (in poly-time)?
- If it is solvable but we don't have a poly time solution, what problem(s) is it most similar too?


## Course Structure

Course divided into three parts:

- Basic automata theory: finite state machines, regular languages, hint of context free languages/grammars, Turing Machines
- Algorithms and algorithm design techniques
- Undecidability and

NP-Completeness, reductions to prove intractability of problems


## Goals

- Algorithmic thinking
- Learn/remember some basic tricks, algorithms, problems, ideas
- Understand/appreciate limits of computation (intractability)
- Appreciate the importance of algorithms in computer science and beyond (engineering, mathematics, natural sciences, social sciences, ...)

Formal languages and complexity (The Blue Weeks!)

## Why Languages?

First 5 weeks devoted to language theory.

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First 5 weeks devoted to language theory.
But why study languages?

## Multiplying Numbers

Consider the following problem:
Problem Given two $n$-digit numbers $x$ and $y$, compute their product.

Grade School Multiplication Compute "partial product" by multiplying each digit of $y$ with $x$ and adding the partial products.

| 3141 |
| ---: |
| $\times 2718$ |
| 25128 |
| 3141 |
| 21987 |
| 6282 |
| 8537238 |



## Time analysis of grade school multiplication

- Each partial product: $\Theta(n)$ time
- Number of partial products: $\leq n$
- Adding partial products: $n$ additions each $\Theta(n)$ (Why?)
- Total time: $\Theta\left(n^{2}\right)$
- Is there a faster way?


## Fast Multiplication

- O( $n^{1.58}$ ) time [Karatsuba 1960] disproving Kolmogorov's belief that $\Omega\left(n^{2}\right)$ is best possible
- O( $n \log n \log \log n)$ [Schonhage-Strassen 1971].

Conjecture: $O(n \log n)$ time possible

- $O\left(n \log n \cdot 2^{O\left(\log ^{*} n\right)}\right)$ time [Furer 2008]
- O( $n \log n$ ) [Harvey-van der Hoeven 2019]

Can we achieve $O(n)$ ? No lower bound beyond trivial one!

## Equivalent Complexity

Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/ QuickSort?

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Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort? How do we compare? The two problems have:

- Different inputs (two numbers vs n-element array)
- Different outputs (a number vs n-element array)
- Different entropy characteristics (from a information theory perspective)


## Languages, Problems and Algorithms ... oh my! II

An algorithm has a runtime complexity.


## Languages, Problems and Algorithms ... oh my! III

A problem has a complexity class!
Recognized by:


Problems do not have run-time since a problem $\neq$ the algorithm used to solve it. Complexity classes are defined differently.
How do we compare problems? What if we just want to know if a problem is "computable".

## Algorithms, Problems and Languages ... oh my! I

Definition

1. An algorithm is a step-by-step way to solve a problem.
2. A problem is some question that we'd like answered given some input. It should be a decision problem of the form "Does a given input fulfill property X."
3. A Language is a set of strings. Given a alphabet, $\Sigma$ a language is a subset of $\Sigma^{*}$

## Algorithms, Problems and Languages ... oh my! I

Definition

1. An algorithm is a step-by-step way to solve a problem.
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3. A Language is a set of strings. Given a alphabet, $\Sigma$ a language is a subset of $\Sigma^{*}$ A language is a formal realization of this problem. For problem X, the corresponding language is:
$L=\{w \mid w$ is the encoding of an input $y$ to problem $X$ and the answer to input y for a problem X is " YES " $\}$
A decision problem $X$ is "YES" is the string is in the language.

## Language of multiplication

How do we define the multiplication problem as a language?
Define $L$ as language where inputs are separated by comma and output is separated by $\mid$.

Machine accepts a $x^{*} y=z$ if " $x^{*} y \mid z$ " is in L. Rejects otherwise.

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$$
\text { Lmult }_{\text {M }}=\left\{\begin{array}{ccc}
1 \times 1 \mid 1, & 1 \times 2 \mid 2, & 1 \times 3 \mid 3, \ldots  \tag{1}\\
2 \times 1 \mid 2, & 2 \times 2 \mid 4, & 2 \times 3 \mid 6, \ldots \\
\vdots & \vdots & \vdots \\
n \times 1 \mid n, & n \times 2 \mid 2 n, & n \times 3 \mid 3 n, \ldots
\end{array}\right\}
$$

## Language of sorting

We do the same thing for sorting.
Define $L$ as language where inputs are separated by comma and output is separated by $\mid$.

Machine accepts a $\left[i_{1}, i_{2}, \ldots\right]=\operatorname{sort}\left(\left\{i_{1}, i_{2}, \ldots\right\}\right)$ if "x[]|z[]" is in $L$. Rejects otherwise.

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$$
L_{\text {sort2 }}=\left\{\begin{array}{ccc}
1,1 \mid 1,1 & 1,2 \mid 1,2 & 1,3 \mid 1,3, \ldots  \tag{2}\\
2,1 \mid 1,2, & 2,2 \mid 2,2, & 2,3 \mid 2,3, \ldots \\
\vdots & \vdots & \vdots \\
n, 1 \mid 1, n, & n, 2 \mid 2, n, & n, 3 \mid 3, n, \ldots
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\end{array}\right\}
$$

If the same type of machine can recognize both languages, then that gives us an upperbound top their hardness.

How do we formulate languages?

## Strings

## Alphabet

An alphabet is a finite set of symbols.
Examples of alphabets:

- $\Sigma=\{0,1\}$,
- $\Sigma=\{a, b, c, \ldots, z\}$,
- ASCII.
- UTF8.
- $\Sigma=$
$\{\langle(\mathrm{w})$ forward $\rangle,\langle(\mathrm{a})$ strafe left $\rangle,\langle(\mathrm{s})$ back $\rangle,\langle(\mathrm{d})$ strafe right $\rangle\}$


## String Definition

## Definition

1．A string／word over $\Sigma$ is a finite sequence of symbols over
$\Sigma$ ．For example，＇0101001＇，＇string＇，＇〈moveback〉〈rotate90〉’
2．$x \cdot y \equiv x y$ is the concatenation of two strings
3．The length of a string $w$（denoted by $|w|$ ）is the number of symbols in $w$ ．For example，$|101|=3,|\epsilon|=0$
4．For integer $n \geq 0, \Sigma^{n}$ is set of all strings over $\Sigma$ of length $n$ ． $\Sigma^{*}$ is the set of all strings over $\Sigma . \mathcal{E}^{n}=\sum_{\varepsilon_{0}} \cdot \sum_{\Sigma_{0}} \cdot \Sigma . . . \varepsilon$
5．$\Sigma^{*}$ set of all strings of all lengths including empty string．
Question：$\left\{{ }^{\prime} 0^{\prime}, 1^{\prime} 1^{\prime}\right\}^{*}=\{0,1,00,01,10,11,000,00,010, \cdots$

## Emptiness

- $\epsilon$ is a string containing no symbols. It is not a set
$\cdot(\{\epsilon\}$ is a set containing one string: the empty string. It is a set, not a string.
$\emptyset \emptyset$ is the empty set. It contains no strings.
Question: What is $|\{\emptyset\}|=1$

$$
\phi=0
$$

## Concatenation and properties

- If $x$ and $y$ are strings then $x y$ denotes their concatenation.
- Concatenation defined recursively:
- $x y=y$ if $x=\epsilon$
- $x y=a(w y)$ if $x=a w$
- $x y$ sometimes written as $x \cdot y$.
- concatenation is associative: $(u v) w=u(v w)$ hence write $u v w \equiv(u v) w=u(v w)$
- not commutative: uv not necessarily equal to vu
- The identity element is the empty string $\epsilon$ :

$$
\epsilon U=U \epsilon=U
$$

## Substrings, prefixes, Suffixes

Definition
$v$ is substring of $w \Longleftrightarrow$ there exist strings $x, y$ such that
$w=x \vee y$.

- If $x=\epsilon$ then $v$ is a prefix of $w \quad w=v y$
- If $y=\epsilon$ then $v$ is a suffix of $w \quad w=x w$


## Subsequence

A subsequence of a string $w[1 \ldots n]$ is either a subsequence of $w[2 \ldots n]$ or $w[1]$ followed by a subsequence of $w[2 \ldots n]$.

## Example <br> EE37 is a subsequence of ECE374B

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## Example

EE37 is a subsequence of ECE374B

Question: How many subsequences are there in a string

$$
|w|=6 ? \quad 2^{6}
$$

$$
\text { a a a } a<a
$$

## String exponent

Definition
If $w$ is a string then $w^{n}$ is defined inductively as follows:
$w^{n}=\epsilon$ if $n=0$
$w \cdot w \cdot$. . . $w$
$w^{n}=w w^{n-1}$ if $n>0$
Question: $(\text { ha })^{3}=$. haha ha

## Rapid-fire questions -strings

Answer the following questions taking $\Sigma=\{0,1\}$.

1. What is $\Sigma^{0}$ ? $\{\varepsilon\}$
2. How many elements are there in $\Sigma^{n}$ ? $2^{n}$
3. If $|u|=2$ and $|v|=3$ then what is $|u \cdot v|$ ? 5
4. Let $u$ be an arbitrary string in $\Sigma^{*}$. What is $\epsilon u$ ? What is $u \epsilon$ ?
u


Languages

$\sum=$ all the string
 nowhere

$$
|\omega|=1
$$

$$
\left\{\begin{array}{l}
0 \\
1
\end{array}\right\}
$$

## Languages

Definition
A language $L$ is a set of strings over $\Sigma^{*}$. In other words $L \subseteq \Sigma^{*}$.

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A language $L$ is a set of strings over $\Sigma$. In other words $L \subseteq \Sigma^{*}$.
Standard set operations apply to languages.

- For languages $A, B$ the concatenation of $A, B$ is $A B=\{x y \mid x \in A, y \in B\}$.
- For languages $A, B$, their union is $A \cup B$, intersection is $A \cap B$, and difference is $A \backslash B$ (also written as $A-B$ ).
- For language $A \subseteq \Sigma^{*}$ the complement of $A$ is $\bar{A}=\Sigma^{*} \backslash A$.


## Set Concatenation

Definition
Given two sets $X$ and $Y$ of strings (over some common alphabet $\Sigma$ ) the concatenation of $X$ and $Y$ is

$$
\begin{equation*}
X Y=\{x y \mid x \in X, y \in Y\} \tag{3}
\end{equation*}
$$

Question: $X=\{E C E, C S 4\}, Y=\{340,374\} \Longrightarrow$
$X Y=$.


## $\Sigma^{*}$ and languages

Definition

$$
\Sigma=\left\{\begin{array}{l}
0 \\
1
\end{array}\right\}^{2}=\left\{\begin{array}{l}
0 \\
1
\end{array}\right\}\left\{\begin{array}{l}
0 \\
1
\end{array}\right\}=\left\{\begin{array}{cc}
00 & 0 \\
10 & 11
\end{array}\right\}
$$

1. $\Sigma^{n}$ is the set of all strings of length $n$. Defined inductively:
2. $\Sigma^{+}=\cup_{n \geq 1} \Sigma^{n}$ is the set of nonempty strings.

## Definition

A language $L$ is a set of strings over $\Sigma$. In other words $L \subseteq \Sigma^{*}$.
Question: Does $\Sigma^{*}$ have strings of infinite length?

Rapid-Fire questions - Languages

Problem
Consider languages over $\Sigma=\{0,1\}$.

1. What is $\emptyset^{0}$ ?
2. If $|L|=2$, then what is $\left|L^{4}\right|$ ?
3. What is $\emptyset^{*},\{\epsilon\}^{*}$ ?
4. For what $L$ is $L^{*}$ finite?
5. What is $\emptyset^{+}$?
6. What is $\{\epsilon\}^{+}$?

## Terminology Review

Let's review what we learned.

- A character $(a, b, c, x)$ is a unit of information represented by a symbol: (letters, digits, whitespace)
- A alphabet $(\Sigma)$ is a set of characters
- A string $(w)$ is a sequence of characters
- A language $(A, B, C, L)$ is a set of strings


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- A string $(w)$ is a sequence of characters
- A language $(A, B, C, L)$ is a set of strings
- A grammar $(G)$ is a set of rules that defines the strings that belong to a language


## Languages: easiest, easy, hard, really hard, really ${ }^{n}$ hard



- Regular languages.
- Regular expressions.
- DFA: Deterministic finite automata.
- NFA: Non-deterministic finite automata.
- Languages that are not regular.
- Context free languages (stack).
- Turing machines: Decidable languages.
- TM Undecidable/unrecognizable languages (halting theorem)


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- Regular languages.
- Regular expressions. $\leftarrow$ Next lecture
- DFA: Deterministic finite automata.
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- Languages that are not regular.
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## That's it for now

Check the course website (https://ecealgo.com) for lab and hw schedule.

