ECE-374-B: Lecture 0 - Logistics and Strings/Languages

Lecturer: Nickvash Kani

August 26, 2025

University of Illinois at Urbana-Champaign

Course Administration

Course Policies

See website

Discussion Sessions/Labs

- 50min problem solving session led by TAs
- Two times a week
- Go to your assigned discussion section
- Bring pen and paper!

No Homeworks ... but there is a catch

- I have lost faith in homeworks being an effective tool for study in the modern age.
- There will be series of short 5 minute quizzes at the beginning of each discussion section.
- Quizzes will gauge mastery of material and serve as a commitment device to study lectures/labs.
- Each quiz is worth 2 points, there are 25 points in your final grade and 20 quizzes during the semester. Seven out of the 20 quizzes can be dropped/failed and still get full credit!
- Quizzes are designed to be easy if you engaged with the lecture and lab material (Quiz 0 will literally be a problem from Lab 0).

Again all policy information should be on course website:
`https://ecealgo.com/`

Any questions?

Over-arching course questions

High-Level Questions

This course introduces three distinct fields of computer science research:

- Computational complexity.
  - Given infinite time and a certain machine, is it possible to solve a given problem.
- Algorithms
  - Given a deterministic Turing machine, how fast can we solve certain problems.
- Limits of computation.
  - Are there tasks that our computers cannot do and how do we identify these problems?

Why not just focus on Algorithms?

When someone asks you, "How fast can you compute problem X", they are actually asking:

- Is X solvable using the deterministic Turing machines we have at our disposal?
- If it is solvable, can we find the solution efficiently (in poly-time)?
- If it is solvable but we don't have a poly time solution, what problem(s) is it most similar too?

Course Structure

Course divided into three parts:

- Basic automata theory: finite state machines, regular languages, hint of context free languages/grammars, Turing Machines
- Algorithms and algorithm design techniques
- Undecidability and NP-Completeness, reductions to prove intractability of problems

Week	Tuesday Lecture	Wed Lab	Thursday Lecture	Fri Lab
Aug 23-26	Adminis trivia and course goals Introduction and history ; strings [Sariel's Videos, Lec 1] [B: slides, scribbles, video]	String induction [Jeff's induction notes , Chandra's induction notes] [solutions][recording]	Languages and regular expressions [Sariel's Videos, Lec 2] [B: slides, scribble]	Regular expressions [solutions][recording]
Aug 30 - Sep 2	DFAs: intuition, definitions, closure properties [Automata Tutor , JFLAP , Mahesh's DFA notes , Sariel's Videos, Lec 3] [B: slides, scribble]	DFA construction [solutions] [recording]	Non-Determinism, NFAs [Sariel's Videos, Lec 4] [B: slides, scribble]	Language transformations [solutions] [recording]
Sep 6-9	Equivalence of DFAs, NFAs, and regular expressions [Sariel's Videos, Lec 5] [B: slides, scribble]	Regex to NFA to DFA (to Regex) [solutions] [recording]	Fooling Sets and Proving Non-Regularity [Mahesh's DFA notes , Fall 2015 TAs' Fooling Sets Notes , Sariel's Videos, Lec 6] [B: slides, scribble]	Proving Non-Regularity [solutions][recording]
Sep 13-16	Context-free languages and grammars [Sariel's Videos, Lec 7] [B: slides, scribble]	Context-free grammars [solutions][recording]	Turing machines: history, formal definitions, examples, variations [Sariel's Videos, Lec 8] [B: slides, scribble]	Turing Machines [solutions][recording]
Sep 20-23	Universal Turing machines [Sariel's Videos, Lec 8] [B: slides, scribble]	Midterm 1 Review (effectively office hours)	**Midterm 1 – Thursday Sep 22**-12:30–13:45 [skill-set , fodder , cheat sheet]	No Instruction
Sep 27-30	Reductions & Recursion [Sariel's Videos, Lec 10 , Notes on Solving Recurrences] [B: slides, scribbles]	Binary search [solutions][recording]	Divide and conquer: Selection, Karatsuba [Sariel's Videos, Lec 11] [B: slides, scribbles]	Divide and Conquer [solutions][recording]
Oct 4-7	Backtracking [Sariel's Videos, Lec 12] [B: slides, scribbles]	Backtracking [solutions][recording]	Dynamic programming [Sariel's Videos, Lec 13] [B: slides, scribbles]	Dynamic programming [solutions][recording]
Oct 11-14	More Dynamic programming [Sariel's Videos, Lec 14] [B: slides, scribbles]	More Dynamic programming [solutions][recording]	Graphs, Basic Search [Chandra's Graph notes , Sariel's Videos, Lec 15] [B: slides, scribbles]	Even more DP [solutions][recording]
Oct 18-21	Directed Graphs, DFS, DAGs and Topological Sort [Chandra's Graph notes , Sariel's Videos, Lec 16] [B: slides, scribbles]	Graph Modeling [solutions][recording]	Shortest Paths: BFS and Dijkstra [Chandra's Graph notes , Sariel's Videos, Lec 17] [B: slides, scribbles]	Shortest Paths [solutions][recording]
Oct 25-28	Bellman-Ford, Dynamic Programming on DAGs [Chandra's Graph notes , Sariel's Videos, Lec 18] [B: slides, scribbles]	More Shortest Paths [solutions][recording]	MST Algorithms [B: slides, scribbles]	MST [solutions][recording]
Nov 1-4	**Midterm 2 (Recursion/DP/Graph Algorithms) – Tuesday, Nov 1 12:30–13:45** [skill-set , fodder , cheat sheet]	No Instruction	Reductions [Sariel's Videos, Lec 21] [B: slides, scribbles]	Reductions [solutions][recording]
Nov 8-11	**SAT, NP and NP-Hardness** **[Sariel's Videos, Lec 22-24]** [B: slides, scribbles]	NP-hardness reductions [solutions][recording]	More NP-Hardness [Sariel's Videos, Lec 23-24] [B: slides, scribbles]	More NP-Hardness [solutions][recording]
Nov 15-18	Undecidability 1 [Sariel's Videos, Lec 9] [B: slides, scribbles]	Undecidability reductions [solutions][recording]	Undecidability 2 [Sariel's Videos, Lec 9] [B: slides, scribbles]	No instruction
	Fall Break (Nov 19-27). Have fun.			
Nov 29 - Dec 2	*Optional review for Midterm 3* [B: slides, scribbles]	*Optional Review for midterm 3 (effectively office hours)*	**Midterm 3 (Reductions/P-NP/Decidability) – Thursday, Dec 1 15:30–16:45** [skill-set , fodder , cheat sheet]	
Dec 6-9	Wrap-up, closing remarks *Optional review for Final Exam* [B: slides, scribbles]	*Optional Review for final exam*	**Reading Day** **ICES Forms Due**	

Final Exam – TBD [Skill set , cheat sheet]

Goals

- Algorithmic thinking
- Learn/remember some basic tricks, algorithms, problems, ideas
- Understand/appreciate limits of computation (intractability)
- Appreciate the importance of algorithms in computer science and beyond (engineering, mathematics, natural sciences, social sciences, ...)

Formal languages and complexity (The Blue Weeks!)

Why Languages?

First 5 weeks devoted to language theory.

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But why study languages?

Multiplying Numbers

Consider the following problem:

> **Problem** Given two n-digit numbers x and y, compute their product.

Grade School Multiplication
Compute "partial product" by multiplying each digit of y with x and adding the partial products.

$$
\begin{array}{r}
3141 \\
\times 2718 \\
\hline
25128 \\
3141 \\
21987 \\
6282 \\
\hline
8537238
\end{array}
$$

$O(n^2)$

Time analysis of grade school multiplication

- Each partial product: $\Theta(n)$ time
- Number of partial products: $\leq n$
- Adding partial products: n additions each $\Theta(n)$ (Why?)
- Total time: $\Theta(n^2)$
- Is there a faster way?

Fast Multiplication

- $O(n^{1.58})$ time [Karatsuba 1960] disproving Kolmogorov's belief that $\Omega(n^2)$ is best possible
- $O(n \log n \log \log n)$ [Schonhage-Strassen 1971].
  Conjecture: $O(n \log n)$ time possible
- $O(n \log n \cdot 2^{O(\log^* n)})$ time [Furer 2008]
- $O(n \log n)$ [Harvey-van der Hoeven 2019]

Can we achieve $O(n)$? No lower bound beyond trivial one!

Equivalent Complexity

Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort?

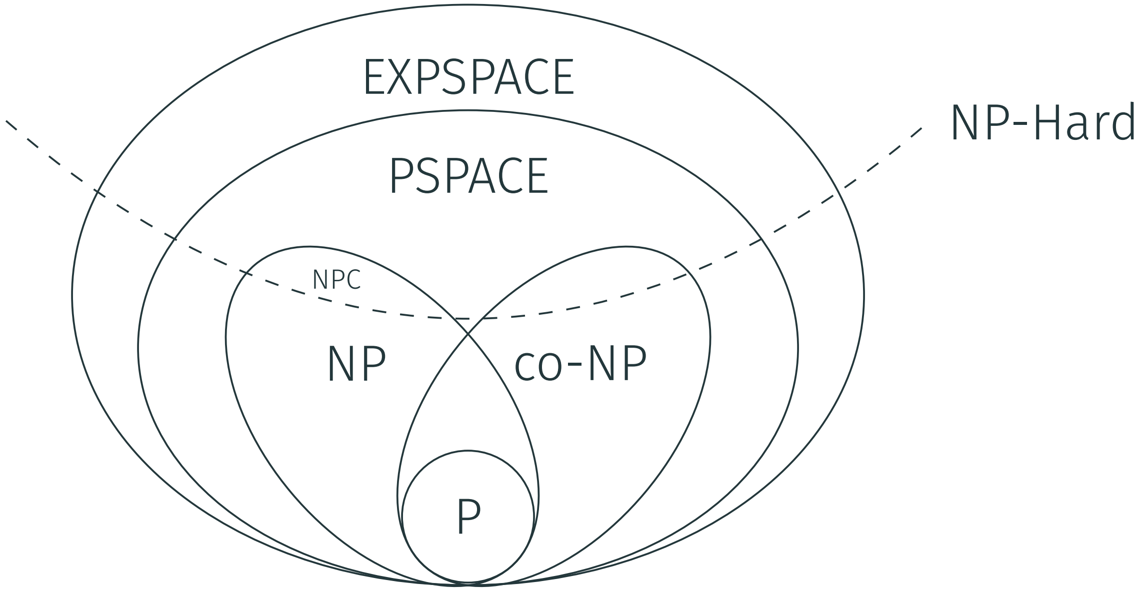
Equivalent Complexity

Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort? How do we compare? The two problems have:

- Different inputs (two numbers vs n-element array)
- Different outputs (a number vs n-element array)
- Different entropy characteristics (from a information theory perspective)

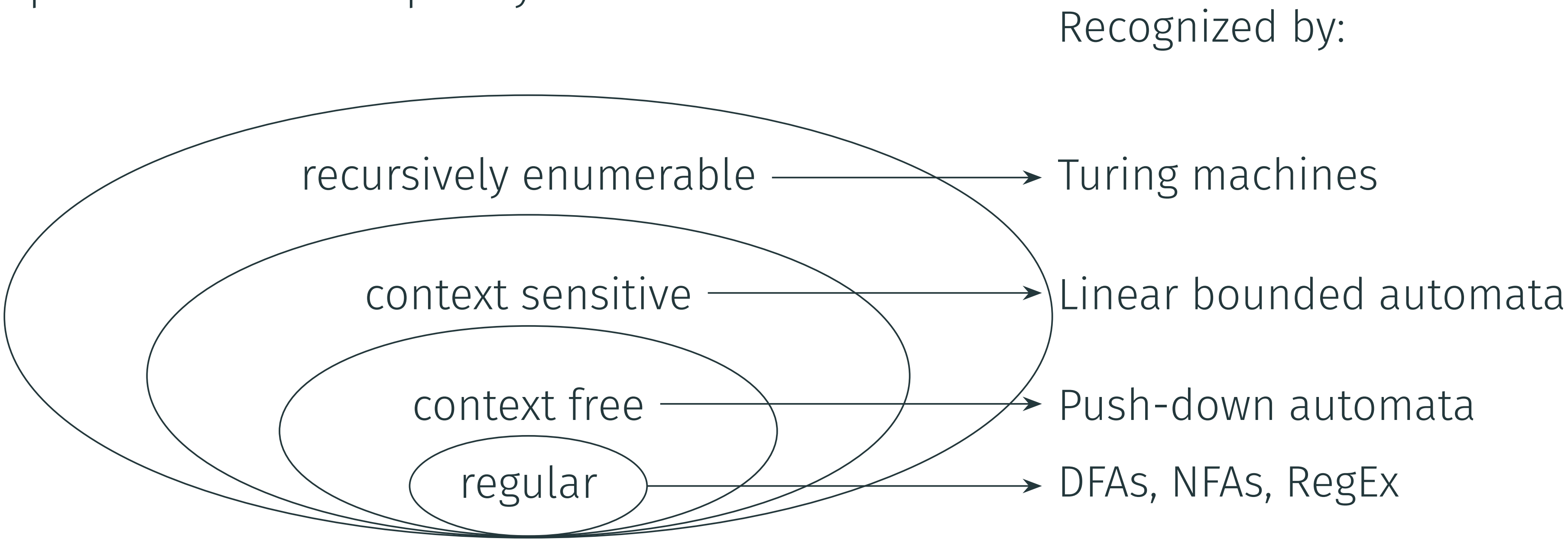
Languages, Problems and Algorithms ... oh my!

An algorithm has a runtime complexity.



Languages, Problems and Algorithms ... oh my!

A problem has a complexity class!



Problems do not have run-time since a problem $\neq$ the algorithm used to solve it.
Complexity classes are defined differently.
How do we compare problems? What if we just want to know if a problem is "computable".

Algorithms, Problems and Languages ... oh my! I

Definition

1. An algorithm is a step-by-step way to solve a problem.
2. A problem is some question that we'd like answered given some input. It should be a decision problem of the form "Does a given input fulfill property X."
3. A Language is a set of strings. Given a alphabet, Σ a language is a subset of Σ^*

Definition

1. An algorithm is a step-by-step way to solve a problem.
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   L = {w | w is the encoding of an input y to problem X and the answer to input y for a problem X is "YES" }
   A decision problem X is "YES" is the string is in the language.

Language of multiplication

How do we define the multiplication problem as a language?

Define L as language where inputs are separated by comma and output is separated by |.

Machine accepts a x*y=z if "x*y|z" is in L. Rejects otherwise.

Problem

What is $x * y \rightarrow z$

Decision Problem

Is $x \cdot y = z$ → Yes / → No

Sorting Decision

Is the sorted version of $A = B$ → Yes / → No

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$$L_{MULT2} = \left\{ \begin{matrix} 1 \times 1|1, & 1 \times 2|2, & 1 \times 3|3, \dots \\ 2 \times 1|2, & 2 \times 2|4, & 2 \times 3|6, \dots \\ \vdots & \vdots & \vdots \\ n \times 1|n, & n \times 2|2n, & n \times 3|3n, \dots \end{matrix} \right\} \tag{1}$$

Language of sorting

We do the same thing for sorting.

Define L as language where inputs are separated by comma and output is separated by |.

Machine accepts a $[i_1, i_2, \ldots] = sort(\{i_1, i_2, \ldots\})$ if ”x[]|z[]” is in L. Rejects otherwise.

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$$L_{Sort2} = \begin{Bmatrix} 1,1|1,1 & 1,2|1,2 & 1,3|1,3,\ldots \\ 2,1|1,2, & 2,2|2,2, & 2,3|2,3,\ldots \\ \vdots & \vdots & \vdots \\ n,1|1,n, & n,2|2,n, & n,3|3,n,\ldots \end{Bmatrix} \tag{2}$$

Language of sorting

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If the same type of machine can recognize both languages, then that gives us an upperbound top their hardness.

How do we formulate languages?

Strings

Alphabet

An alphabet is a **finite** set of symbols.

Examples of alphabets:

- $\Sigma = \{0, 1\}$,
- $\Sigma = \{a, b, c, \ldots, z\}$,
- ASCII.
- UTF8.
- $\Sigma = \{\langle\text{(w)forward}\rangle, \ \langle\text{(a)strafe left}\rangle, \ \langle\text{(s)back}\rangle, \ \langle\text{(d)strafe right}\rangle\}$

String Definition

Definition

1. A string/word over Σ is a **finite sequence** of symbols over Σ. For example, '0101001', '*string*', '$\langle\text{moveback}\rangle\langle\text{rotate90}\rangle$'
2. $x \cdot y \equiv xy$ is the concatenation of two strings
3. The length of a string w (denoted by $|w|$) is the number of symbols in w. For example, $|101| = 3$, $|\epsilon| = 0$
4. For integer $n \geq 0$, Σ^n is set of all strings over Σ of length n. Σ^* is the set of all strings over Σ. $\Sigma = \{0, 1\}$ $\Sigma^2 = \{00, 01, 10, 11\}$
5. Σ^* set of all strings of all lengths including empty string. $\Sigma \cdot \Sigma$

* represent all positive #s

Question: $\{'0', '1'\}^* =$

$\Sigma^0 = \epsilon$ $\{\epsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, 011, \ldots\}$ $\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \Sigma^3 \cup \Sigma^4 \cup \ldots$ $\{0,1\} \cdot \{0,1\}$

Emptiness

- ϵ is a string containing no symbols. It is not a set
- $1 = |\{\epsilon\}|$ is a set containing one string: the empty string. It is a set, not a string.
- $\emptyset$ is the empty set. It contains no strings.

Question: What is $|\{\emptyset\}| = 1$

sets that contains the empty set

$|\emptyset| = 0$

Concatenation and properties

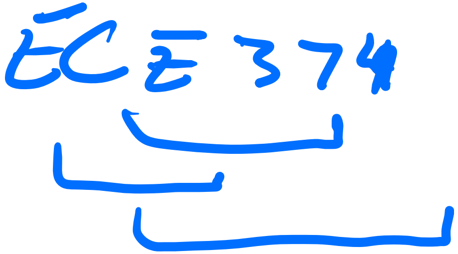
- If x and y are strings then xy denotes their concatenation.
- Concatenation defined recursively :
  - $xy = y$ if $x = \epsilon$
  - $xy = a(wy)$ if $x = aw$
- xy sometimes written as $x \bullet y$.
- concatenation is associative: $(uv)w = u(vw)$ hence write $uvw \equiv (uv)w = u(vw)$
- **not** commutative: uv not necessarily equal to vu
- The identity element is the empty string ϵ:

$$\epsilon u = u\epsilon = u.$$

Substrings, prefixes, Suffixes

Definition
v is substring of $w \iff$ there exist strings x, y such that $w = xvy$.

- If $x = \epsilon$ then v is a prefix of w
- If $y = \epsilon$ then v is a suffix of w



Subsequence

A subsequence of a string $w[1...n]$ is either a subsequence of $w[2...n]$ or $w[1]$ followed by a subsequence of $w[2...n]$.

Example

EE37 is a subsequence of *ECE374B*

Subsequence

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Example

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Question: How many sub-sequences are there in a string $|w| = 6$?

$w =$ _ _ _ _ _ _

2^n

String exponent

Definition
If w is a string then w^n is defined inductively as follows:

$w^n = \epsilon$ if $n = 0$

$w^n = ww^{n-1}$ if $n > 0$

Question: $(ha)^3 =$. hahaha

Rapid-fire questions -strings

Answer the following questions taking $\Sigma = \{0, 1\}$.

1. What is Σ^0?
2. How many elements are there in Σ^n?
3. If $|u| = 2$ and $|v| = 3$ then what is $|u \bullet v|$?
4. Let u be an arbitrary string in Σ^*. What is ϵu? What is $u\epsilon$?

Languages

Languages

Definition
A language L is a set of strings over Σ. In other words $L \subseteq \Sigma^*$.

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Standard set operations apply to languages.

- For languages A, B the concatenation of A, B is $AB = \{xy \mid x \in A, y \in B\}$.
- For languages A, B, their union is $A \cup B$, intersection is $A \cap B$, and difference is $A \setminus B$ (also written as $A - B$).
- For language $A \subseteq \Sigma^*$ the complement of A is $\bar{A} = \Sigma^* \setminus A$.

Set Concatenation

Definition
Given two sets X and Y of strings (over some common alphabet Σ) the concatenation of X and Y is

$$XY = \{xy \mid x \in X, y \in Y\} \tag{3}$$

Question: $X = \{ECE, CS,\}, Y = \{340, 374\} \implies$
$XY = .$

Definition

1. Σ^n is the set of all strings of length n. Defined inductively:
   $\Sigma^n = \{\epsilon\}$ if $n = 0$
   $\Sigma^n = \Sigma\Sigma^{n-1}$ if $n > 0$
2. $\Sigma^* = \cup_{n\geq 0}\Sigma^n$ is the set of all finite length strings
3. $\Sigma^+ = \cup_{n\geq 1}\Sigma^n$ is the set of non-empty strings.

Definition

A language L is a set of strings over Σ. In other words $L \subseteq \Sigma^*$.

Question: Does Σ^* have strings of infinite length?

Rapid-Fire questions - Languages

Problem
Consider languages over $\Sigma = \{0, 1\}$.

1. *What is* $\emptyset^0$?
2. *If* $|L| = 2$, *then what is* $|L^4|$?
3. *What is* $\emptyset^*$, $\{\epsilon\}^*$?
4. *For what* L *is* L^* *finite?*
5. *What is* $\emptyset^+$?
6. *What is* $\{\epsilon\}^+$?

Terminology Review

Let's review what we learned.

- A character(a, b, c, x) is a unit of information represented by a symbol: (letters, digits, whitespace)
- A alphabet(Σ) is a set of characters
- A string(w) is a sequence of characters
- A language(A, B, C, L) is a set of strings

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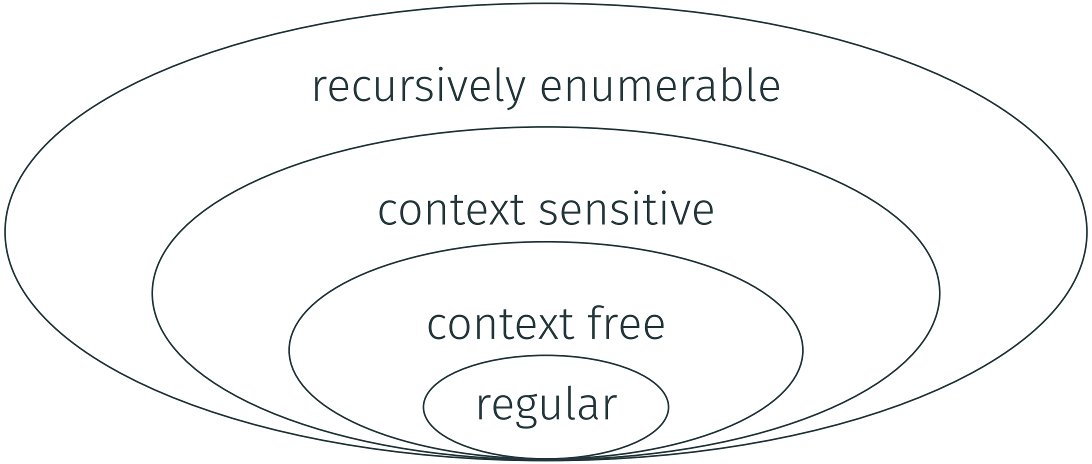
- A string(w) is a sequence of characters
- A language(A, B, C, L) is a set of strings
- A grammar(G) is a set of rules that defines the strings that belong to a language

Languages: easiest, easy, hard, really hard, really^n hard



- Regular languages.
    - Regular expressions.
    - DFA: Deterministic finite automata.
    - NFA: Non-deterministic finite automata.
- Context free languages (stack).
- Turing machines: Decidable languages.
- TM Undecidable/unrecognizable languages (halting theorem).

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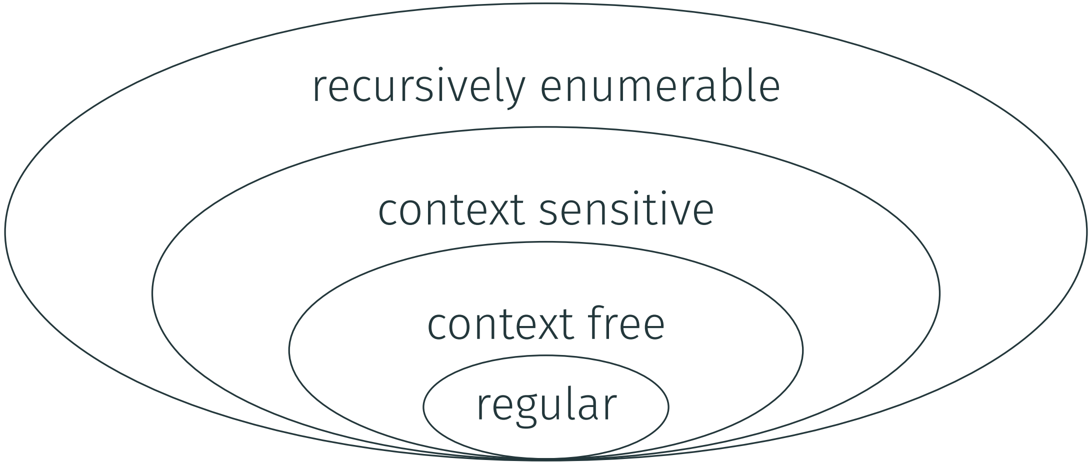

- Regular languages.
    - Regular expressions. $\leftarrow$ **Next lecture**
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That’s it for now

Check the course website (`https://ecealgo.com/`) for course schedule(s).

OH will begin Thursday.