

# ECE 374 B ✧ Spring 2024

## 🌀 Homework 4 🌀

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- **Submit your solutions electronically on the course Gradescope site as PDF files.** If you plan to typeset your solutions, please use the  $\text{\LaTeX}$  solution template on the course web site. If you must submit scanned handwritten solutions, please use a black pen on blank white paper and a high-quality scanner app (or an actual scanner, not just a phone camera).
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### 👉 **Some important course policies** 👈

- **You may use any source at your disposal**—paper, electronic, or human—but you *must* cite *every* source that you use, and you *must* write everything yourself in your own words. See the academic integrity policies on the course web site for more details.
  - **Avoid the Three Deadly Sins!** Any homework or exam solution that breaks any of the following rules will be given an *automatic zero*, unless the solution is otherwise perfect. Yes, we really mean it. We're not trying to be scary or petty (Honest!), but we do want to break a few common bad habits that seriously impede mastery of the course material.
    - Always give complete solutions, not just examples.
    - Always declare all your variables, in English. In particular, always describe the specific problem your algorithm is supposed to solve.
    - Never use weak induction.
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### **See the course web site for more information.**

If you have any questions about these policies, please don't hesitate to ask in class, in office hours, or on Piazza.

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1. Solve the following recurrence relations. For parts (a) and (b), provide an exact solution. For parts (c) and (d), provide an asymptotic upper bound. For both cases, your solution must explain how you obtained the expression.

(a)  $A(n) = A(n-1) + 2n + 1; A(0) = 0$

(b)  $B(n) = B(n-1) + n(n-1) - 1; B(0) = 0$

(c)  $C(n) = C(n/2) + C(n/3) + C(n/6) + n$

(d)  $D(n) = D(n/2) + D(n/3) + D(n/6) + n^2$

2. Consider the following variants of the Towers of Hanoi. For each of variant, describe an algorithm to solve it in as few moves as possible. Prove that your algorithm is correct.

Initially, all the  $n$  disks are on peg 1, and you need to move the disks to peg 2. In all the following variants, you are not allowed to put a bigger disk on top of a smaller disk.

- (a) **Hanoi 1:** Suppose you are forbidden to move any disk directly between peg 1 and peg 2, and every move must involve (the third peg) 0. Exactly (i.e., not asymptotically) how many moves does your algorithm make as a function of  $n$ ?

- (b) **Hanoi 2:** Suppose you are only allowed to move disks from peg 0 to peg 1, from peg 1 to peg 2, or from peg 2 to peg 0.

Provide an upper bound, as tight as possible, on the number of moves that your algorithm uses.

(One can derive the exact upper bound by solving the recurrence, but this is too tedious and not required here.)

- (c) **Hanoi 3:** Finally consider the disappearing Tower of Hanoi puzzle where the largest remaining disk will disappear if there is nothing on top of it. The goal here is to get all the disks to disappear and be left with three empty pegs (in as few moves as possible).

Provide an upper bound, as tight as possible, on the number of moves your algorithm uses.

