

ECE-374-B: Lecture 0 - Logistics and Strings/Languages

Instructor: Abhishek Kumar Umrawal

January 16, 2024

University of Illinois at Urbana-Champaign

Course Administration

Instructional Staff

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- **Office hours:** TBD, See course webpage
- **Contacting us:** Use private notes on Piazza to reach course staff. Direct email only for sensitive or confidential information.

About your instructor – Basic info

- **Name:** Abhishek Kumar Umrawal
- **Webpage:** ece.illinois.edu/about/directory/faculty/aumrawal
- **Email:** aumrawal@illinois.edu
- **Office:** ECEB 3054
- **Office hours:** TBD

About your instructor – Education

- **Purdue University**, Ph.D. in Industrial Engineering
Dissertation: Machine Learning Algorithms for Influence Maximization on Social Networks
- **Purdue University**, MS in Economics
- **Indian Institute of Technology (IIT) Kanpur**, MS in Statistics

About your instructor – Prior teaching experience

- **University of Maryland**, Visiting Lecturer of Computer Science and Electrical Engineering

About your instructor – Research interests

Core areas:

1. Combinatorial optimization
2. Approximation algorithms
3. Statistical learning theory
4. Reinforcement learning (RL)
5. Causal inference

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2. Promotional marketing
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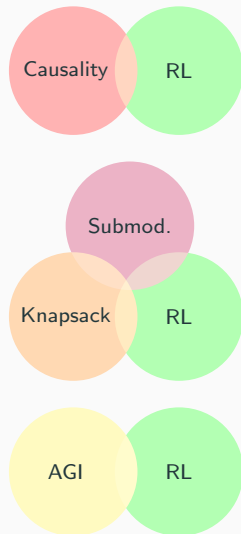
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Intersections:



Working with me on research

If you are interested in working with me then **please send me an email** with subject line 'Expressing interest in working with you on research' **with a brief description of your interests and skills** with no attachments. Please do so only after the second midterm so that you have spent enough time learning algorithms.

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Preferred (but not required) **skills**:

- Mathematical thinking
- Probability and statistics
- Python programming – graphs, object-oriented programming, recursion, etc.
- Algorithms (you're doing it this semester!)

You may fill out [this](#) form to provide further information.

Section A vs B

This semester, the two sections will be run completely **independently**.

- Different lectures.
- Different homeworks, quizzes, exams.
- Different grading policies.

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- Different grading policies.

Section B will be in-person only. Recordings will be attempted but not guaranteed.

Online resources

- **Webpage:** General information, announcements, homeworks, quizzes, course policies will be available at <https://ecealgo.com>.
- **Submission(Gradescope):** Written homework submission and grading, regrade requests. Exams will be uploaded there as well.
- **Communication(Piazza):** Announcements, online questions and discussion, contacting course staff (via private notes).
- **Gradebook (Canvas):** Announcements, online questions and discussion, contacting course staff (via private notes).

See course webpage for links.

Important: Check Piazza/course web page at least once each day.

Discussion Sessions/Labs

- 50min problem solving session led by TAs.
- Two times a week.
- Go to your assigned discussion section.
- Bring pen and paper!

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Discussion sections will have questions that appear on the homework. If, you skip, you're just making more work for yourself later.

Any questions

Again all policy information should be on course website:

<https://ecealgo.com>.

Any questions?

Over-arching course questions

High-Level Questions

This course introduces three distinct fields of computer science research:

- Computational complexity.
 - Given infinite time and a certain machine, is it possible to solve a given problem.
- Algorithms.
 - Given a deterministic Turing machine, how fast can we solve certain problems.
- Limits of computation.
 - Are there tasks that our computers cannot do and how do we identify these problems?

Why not just focus on Algorithms?

When someone asks you, “How fast can you compute problem X ”, they are actually asking:

- Is X solvable using the deterministic Turing machines we have at our disposal?
- If it is solvable, can we find the solution efficiently (in poly-time)?
- If it is solvable but we don't have a poly time solution, what problem(s) is it most similar too?

Course Structure

Course divided into three parts:

- Basic automata theory: finite state machines, regular languages, hint of context free languages/grammars, Turing Machines.
- Algorithms and algorithm design techniques.
- Undecidability and NP-Completeness, reductions to prove intractability of problems.

Week	Thematic Lecture	Workshop	Thematic Lecture	Self-Learn
Aug 26	Introduction and course goals Introduction and History J. Aho & Sethi (1987), p. 1-11 30 mins, written	Defining Automata J. Aho & Sethi (1987), p. 12-13 30 mins, written	Languages and regular expressions Sethi & Aho (1987), p. 14-15 30 mins, written	Regular automata & Sethi & Aho (1987), p. 16-17 30 mins, written
Aug 30 - Sep 2	Finite Automata, Regular Expressions Hopcroft & Ullman (1979), p. 1-11 30 mins, written	FA, NFA, DFA Hopcroft & Ullman (1979), p. 12-13 30 mins, written	Non-deterministic FA, NFA Hopcroft & Ullman (1979), p. 14-15 30 mins, written	Language characterization & Sethi & Aho (1987), p. 16-17 30 mins, written
Sep 6	Existence of DFA, NFA, and regular expressions Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Regex to DFA, DFA to Regex Hopcroft & Ullman (1979), p. 12-13 30 mins, written	Minimal DFA Hopcroft & Ullman (1979), p. 14-15 30 mins, written	Regex Minimization & Sethi & Aho (1987), p. 16-17 30 mins, written
Sep 13	Context free languages and automata Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Context free languages Hopcroft & Ullman (1979), p. 12-13 30 mins, written	Pushdown Automata, Turing Machine Hopcroft & Ullman (1979), p. 14-15 30 mins, written	Turing Machines & Hopcroft & Ullman (1979), p. 16-17 30 mins, written
Sep 20	Context free languages, Turing Machines Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Minimal Turing Machine Hopcroft & Ullman (1979), p. 12-13 30 mins, written	Minimal Turing Machine Hopcroft & Ullman (1979), p. 14-15 30 mins, written	No instruction
Sep 27	Reductions & Decision Problems Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Many-one, & Sethi & Aho (1987), p. 12-13 30 mins, written	Diagonalization, Selection, Recursion Sethi & Aho (1987), p. 14-15 30 mins, written	Diagonalization & Sethi & Aho (1987), p. 16-17 30 mins, written
Oct 4	Reductions & Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Reductions & Sethi & Aho (1987), p. 12-13 30 mins, written	Diagonalization & Sethi & Aho (1987), p. 14-15 30 mins, written	Diagonalization & Sethi & Aho (1987), p. 16-17 30 mins, written
Oct 11	More Decision problems Hopcroft & Ullman (1979), p. 1-11 30 mins, written	More Decision problems Sethi & Aho (1987), p. 12-13 30 mins, written	Complexity, Cook's Theorem Sethi & Aho (1987), p. 14-15 30 mins, written	Complexity & Sethi & Aho (1987), p. 16-17 30 mins, written
Oct 18	Complexity, P, NP, Co-NP and Diagonalization Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Complexity, P, NP, Co-NP Hopcroft & Ullman (1979), p. 12-13 30 mins, written	Complexity, Cook's Theorem Sethi & Aho (1987), p. 14-15 30 mins, written	Complexity & Sethi & Aho (1987), p. 16-17 30 mins, written
Oct 25	Reductions, Cook's Theorem, Complexity Hopcroft & Ullman (1979), p. 1-11 30 mins, written	More Decision Problems Hopcroft & Ullman (1979), p. 12-13 30 mins, written	NP-Completeness Hopcroft & Ullman (1979), p. 14-15 30 mins, written	NP-Completeness & Hopcroft & Ullman (1979), p. 16-17 30 mins, written
Nov 1	NP-Completeness, Reductions Hopcroft & Ullman (1979), p. 1-11 30 mins, written	No instruction	NP-Completeness Hopcroft & Ullman (1979), p. 12-13 30 mins, written	NP-Completeness & Hopcroft & Ullman (1979), p. 16-17 30 mins, written
Nov 8	NP-Completeness, Reductions Hopcroft & Ullman (1979), p. 1-11 30 mins, written	NP-Completeness Hopcroft & Ullman (1979), p. 12-13 30 mins, written	NP-Completeness Hopcroft & Ullman (1979), p. 14-15 30 mins, written	NP-Completeness & Hopcroft & Ullman (1979), p. 16-17 30 mins, written
Nov 15	Undecidability Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Undecidability Hopcroft & Ullman (1979), p. 12-13 30 mins, written	Undecidability Hopcroft & Ullman (1979), p. 14-15 30 mins, written	No instruction
Full Break (Nov 19-21), Exam Week				
Nov 29 - Dec 2	Optional Review for Midterm 1 Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Optional Review for Midterm 1 Hopcroft & Ullman (1979), p. 12-13 30 mins, written	Optional Review for Midterm 1 Hopcroft & Ullman (1979), p. 14-15 30 mins, written	Optional Review for Midterm 1 Hopcroft & Ullman (1979), p. 16-17 30 mins, written
Dec 9	Optional Review for Final Exam Hopcroft & Ullman (1979), p. 1-11 30 mins, written	Optional Review for Final Exam Hopcroft & Ullman (1979), p. 12-13 30 mins, written	Optional Review for Final Exam Hopcroft & Ullman (1979), p. 14-15 30 mins, written	Optional Review for Final Exam Hopcroft & Ullman (1979), p. 16-17 30 mins, written
Final Exam - 180 (300) and 1, cheat sheet & 1				

Goals

- Algorithmic thinking.
- Learn/remember some basic tricks, algorithms, problems, ideas.
- Understand/appreciate limits of computation (intractability).
- Appreciate the importance of algorithms in computer science and beyond (engineering, mathematics, natural sciences, social sciences, ...).

Formal languages and complexity (The Blue Weeks!)

Why Languages?

First 5 weeks devoted to language theory.

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First 5 weeks devoted to language theory.

But why study languages?

Multiplying Numbers

Consider the following problem:

Problem Given two n -digit numbers x and y , compute their product.

Grade School Multiplication

Compute “partial product” by multiplying each digit of y with x and adding the partial products.

$$\begin{array}{r} 3141 \\ \times 2718 \\ \hline 25128 \\ 3141 \\ 21987 \\ 6282 \\ \hline 8537238 \end{array}$$

Time analysis of grade school multiplication

- Each partial product: $\Theta(n)$ time
- Number of partial products: $\leq n$
- Adding partial products: n additions each $\Theta(n)$ (Why?)
- Total time: $\Theta(n^2)$
- Is there a faster way?

Fast Multiplication

- $O(n^{1.58})$ time [Karatsuba 1960] disproving Kolmogorov's belief that $\Omega(n^2)$ is best possible.
- $O(n \log n \log \log n)$ [Schönhage-Strassen. 1971].
Conjecture: $O(n \log n)$ time possible.
- $O(n \log n \cdot 2^{O(\log^* n)})$ time [Furer 2008].
- $O(n \log n)$ [Harvey-van der Hoeven 2019].

Can we achieve $O(n)$? No lower bound beyond trivial one!

Equivalent Complexity

Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort?

Equivalent Complexity

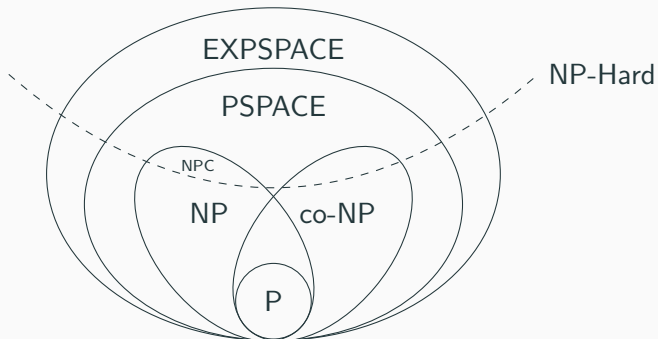
Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort?

How do we compare? The two problems have:

- Different inputs (two numbers vs n-element array).
- Different outputs (a number vs n-element array).
- Different entropy characteristics (from an information theory perspective).

Languages, Problems and Algorithms ... oh my! II

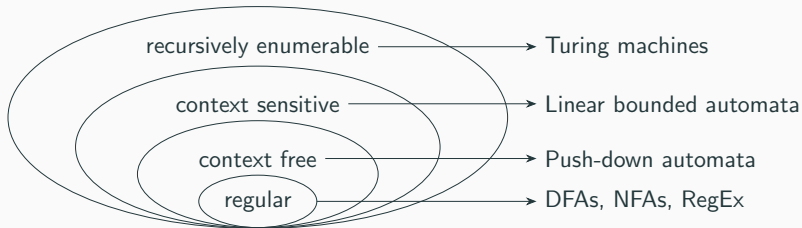
An algorithm has a runtime complexity.



Languages, Problems and Algorithms ... oh my! III

A problem has a complexity class!

Recognized by:



Problems do not have run-time since a problem \neq the algorithm used to solve it. *Complexity classes are defined differently.*

How do we compare problems? What if we just want to know if a problem is “computable”.

Definition

1. An **algorithm** is a step-by-step way to solve a problem.
2. A **problem** is some question that we'd like answered given some input. It should be a decision problem of the form "Does a given input fulfill property X."
3. A **Language** is a set of strings. Given an alphabet, Σ a language is a subset of Σ^* .

Definition

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3. A **Language** is a set of strings. Given a alphabet, Σ a language is a subset of Σ^* . A language is a formal realization of this problem. For problem X, the corresponding language is:

$L = \{w \mid w \text{ is the encoding of an input } y \text{ to problem } X \text{ and the answer to input } y \text{ for a problem } X \text{ is "YES"} \}$

A decision problem X is "YES" if the string is in the language.

Language of multiplication

How do we define the multiplication problem as a language?

Define L as language where inputs are separated by comma and output is separated by $—$.

Machine accepts a $x*y=z$ if " $x*y—z$ " is in L . Rejects otherwise.

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How do we define the multiplication problem as a language?

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Machine accepts a $x*y=z$ if " $x*y—z$ " is in L. Rejects otherwise.

$$L_{MULT2} = \left\{ \begin{array}{lll} 1 \times 1|1, & 1 \times 2|2, & 1 \times 3|3, \dots \\ 2 \times 1|2, & 2 \times 2|4, & 2 \times 3|6, \dots \\ \vdots & \vdots & \vdots \\ n \times 1|n, & n \times 2|2n, & n \times 3|3n, \dots \end{array} \right\} \quad (1)$$

Language of sorting

We do the same thing for sorting.

Define L as language where inputs are separated by comma and output is separated by —.

Machine accepts a $[i_1, i_2, \dots] = \text{sort}(\{i_1, i_2, \dots\})$ if " $x[]—z[]$ " is in L. Rejects otherwise.

Language of sorting

We do the same thing for sorting.

Define L as language where inputs are separated by comma and output is separated by $|$.

Machine accepts a $[i_1, i_2, \dots] = \text{sort}(\{i_1, i_2, \dots\})$ if " $x|z$ " is in L . Rejects otherwise.

$$L_{\text{Sort2}} = \left\{ \begin{array}{ccc} 1, 1|1, 1 & 1, 2|1, 2 & 1, 3|1, 3, \dots \\ 2, 1|1, 2, & 2, 2|2, 2, & 2, 3|2, 3, \dots \\ \vdots & \vdots & \vdots \\ n, 1|1, n, & n, 2|2, n, & n, 3|3, n, \dots \end{array} \right\} \quad (2)$$

Language of sorting

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If the same type of machine can recognize both languages, then that gives us an upperbound to their hardness.

How do we formulate languages?

Strings

Alphabet

An **alphabet** is a **finite** set of symbols.

Examples of alphabets:

- $\Sigma = \{0, 1\}$,
- $\Sigma = \{a, b, c, \dots, z\}$,
- ASCII.
- UTF8.
- $\Sigma = \{\langle(w)forward\rangle, \langle(a)strafe\ left\rangle, \langle(s)back\rangle, \langle(d)strafe\ right\rangle\}$

String Definition

Definition

1. A **string/word** over Σ is a **finite sequence** of symbols over Σ .
For example, '0101001', '*string*', ' $\langle \text{moveback} \rangle \langle \text{rotate90} \rangle$ '
2. $x \cdot y \equiv xy$ is the concatenation of two strings
3. The **length** of a string w (denoted by $|w|$) is the number of symbols in w . For example, $|101| = 3$, $|\epsilon| = 0$
4. For integer $n \geq 0$, Σ^n is set of all strings over Σ of length n .
 Σ^* is the set of all strings over Σ .
5. Σ^* set of all strings of all lengths including empty string.

Question: What is $\{ '0', '1' \}^*$?

Emptiness

- ϵ is a **string** containing no symbols. It is not a set
- $\{\epsilon\}$ is a **set** containing one string: the empty string. It is a set, not a string.
- \emptyset is the **empty set**. It contains no strings.

Question: What is $\{\emptyset\}$?

Concatenation and properties

- If x and y are strings then xy denotes their concatenation.
- **Concatenation** defined recursively :
 - $xy = y$ if $x = \epsilon$
 - $xy = a(wy)$ if $x = aw$
- xy sometimes written as $x \cdot y$.
- concatenation is **associative**: $(uv)w = u(vw)$ hence write $uvw \equiv (uv)w = u(vw)$
- **not** commutative: uv not necessarily equal to vu
- The identity element is the empty string ϵ :

$$\epsilon u = u \epsilon = u.$$

Definition

v is **substring** of $w \iff$ there exist strings x, y such that $w = xvy$.

- If $x = \epsilon$ then v is a **prefix** of w
- If $y = \epsilon$ then v is a **suffix** of w

Subsequence

A subsequence of a string $w[1\dots n]$ is either a subsequence of $w[2\dots n]$ or $w[1]$ followed by a subsequence of $w[2\dots n]$.

Example

EE37 is a subsequence of *ECE374B*

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A subsequence of a string $w[1\dots n]$ is either a subsequence of $w[2\dots n]$ or $w[1]$ followed by a subsequence of $w[2\dots n]$.

Example

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Question: How many sub-sequences are there in a string $|w| = 6$?

Definition

If w is a string then w^n is defined inductively as follows:

$$w^n = \epsilon \text{ if } n = 0$$

$$w^n = ww^{n-1} \text{ if } n > 0$$

Question: $(ha)^3 =$.

Rapid-fire questions -strings

Answer the following questions taking $\Sigma = \{0, 1\}$.

1. What is Σ^0 ?
2. How many elements are there in Σ^n ?
3. If $|u| = 2$ and $|v| = 3$ then what is $|u \cdot v|$?
4. Let u be an arbitrary string in Σ^* . What is ϵu ? What is $u \epsilon$?

Languages

Definition

A **language** L is a set of strings over Σ . In other words $L \subseteq \Sigma^*$.

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Standard set operations apply to languages.

- For languages A, B the **concatenation** of A, B is $AB = \{xy \mid x \in A, y \in B\}$.
- For languages A, B , their **union** is $A \cup B$, **intersection** is $A \cap B$, and **difference** is $A \setminus B$ (also written as $A - B$).
- For language $A \subseteq \Sigma^*$ the **complement** of A is $\bar{A} = \Sigma^* \setminus A$.

Set Concatenation

Definition

Given two sets X and Y of strings (over some common alphabet Σ) the **concatenation** of X and Y is

$$XY = \{xy \mid x \in X, y \in Y\} \quad (3)$$

Question: $X = \{ECE, CS, \}$, $Y = \{340, 374\} \implies$
 $XY = .$

Definition

1. Σ^n is the set of all strings of length n . Defined inductively:

$$\Sigma^n = \{\epsilon\} \text{ if } n = 0$$

$$\Sigma^n = \Sigma\Sigma^{n-1} \text{ if } n > 0$$

2. $\Sigma^* = \cup_{n \geq 0} \Sigma^n$ is the set of all finite length strings

3. $\Sigma^+ = \cup_{n \geq 1} \Sigma^n$ is the set of non-empty strings.

Definition

A **language** L is a set of strings over Σ . In other words $L \subseteq \Sigma^*$.

Question: Does Σ^* have strings of infinite length?

Problem

Consider languages over $\Sigma = \{0, 1\}$.

1. What is \emptyset^0 ?
2. If $|L| = 2$, then what is $|L^4|$?
3. What is \emptyset^* , $\{\epsilon\}^*$?
4. For what L is L^* finite?
5. What is \emptyset^+ ?
6. What is $\{\epsilon\}^+$?

Terminology Review

Let's review what we learned.

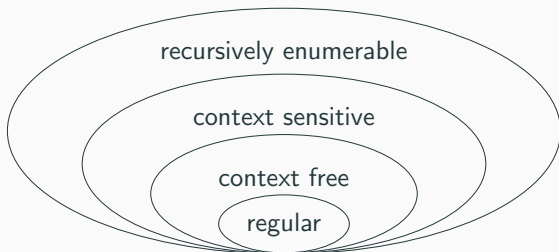
- A **character**(a, b, c, x) is a unit of information represented by a symbol: (letters, digits, whitespace)
- A **alphabet**(Σ) is a set of characters
- A **string**(w) is a sequence of characters
- A **language**(A, B, C, L) is a set of strings

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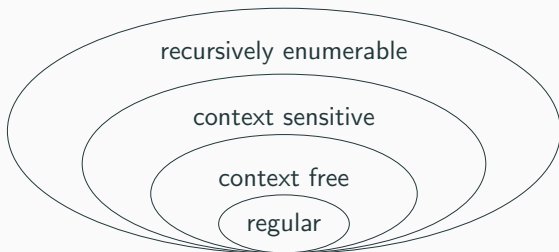
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- A **string**(w) is a sequence of characters
- A **language**(A, B, C, L) is a set of strings
- A **grammar**(G) is a set of rules that defines the strings that belong to a language

Languages: easiest, easy, hard, really hard, reallyⁿ hard



- Regular languages.
 - Regular expressions.
 - DFA: Deterministic finite automata.
 - NFA: Non-deterministic finite automata.
 - Languages that are not regular.
- Context free languages (stack).
- Turing machines: Decidable languages.
- TM Undecidable/unrecognizable languages (halting theorem). 38

Languages: easiest, easy, hard, really hard, reallyⁿ hard



- Regular languages.
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 - DFA: Deterministic finite automata.
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That's it for now

Check the course website (<https://ecealgo.com>) for lab and hw schedule.